## IBM System/360 Model 44

## Programming System

## Assembler Language

This publication contains specifications for the IBM System/360 Nodel 44 Programming System Assembler Language.

This assembler language is used to write programs for the Model 44. The IBM System/360 Model 44 Programming System Assembler program processes the language and provides auxiliary functions useful in the preparation and documentation of a program.

This publication is a reference manual for the programmer using the assembler language and its features.

This publication presents information common to all parts of the language, followed by specific information concerning the symbolic machine instruction codes and the assembler program functions provided for the programmer's use.

Appendixes A through $F$ present such items as a summary chart, for constants, instruction listings, character set representations, and other aids to programming. Appendix $G$ is a features comparison chart of System/360 assemblers.

Knowledge of IBM System/360 machine operations, particularly storage addressing, data formats, and machine instruction formats and functions, is prerequisite to using this publication. It is assumed that the reader has experience with programming concepts and techniques or has completed
basic courses of instruction in these areas.

The publications most closely supplemental to this one are:

IBM System/360: Principles of operation, Form A22-6821

IBM System/360: System Summary, Form A22-6810

IBM System/360 Model 44: Functional Characteristics, Form A22-6875

IBM System 360 Model 44 Programming System: Concepts and Facilities, Form C28-6810

IBM System/360 Model 44 Programming System: Guide to System Use, Form C28-6812

Data Acquisition Special Features for IBM System/ 360 Model 44, Form A22-6900

Second Edition
This is a major revision of, and makes obsolete, c28-6811-0. Section 7, "Update Feature," has been substantially revised to include additional information and examples of update operations. Appendix $G$, "Features comparison Checklists," has been rewritten to define more specifically the relationship between the IBM System 360 Model 44 Programming System Assembler Language and the other System/360 programming support system assembler languages. Because of a change in specifications, all references to the length attribute of a symbol have been deleted. Changes to the text other than these are indicated by a vertical line to the left of the change.

Specifications contained herein are subject to change from time to time. Any such change will be reported in subsequent revisions or Technical Newsletters.

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A form is provided at the back of this publication for reader's comments. If the form has been removed, comments may be addressed to IBM Corporation, Programming Publications, 1271 Avenue of the Americas, New York, N.Y., 10020.

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Computer programs may be expressed in machine language, i.e., language interpreted directly by the computer, or in a symbolic language, which is much more meaningful to the programmer The symbolic language, however, must be translated into machine language before the computer can execute the program. This function is accomplished by a processing program.

Of the various symbolic programming languages, assembler languages are closest to machine language in form and content. The assembler language discussed in this publication is a symbolic programming language for the IBM System/ 360 Model 44. It enables the programmer to use IBM System 360 machine functions as though he were coding in System/360 Model 44 machine language.

The assembler program that processes the language translates symbolic instructions into machine-language instructions, assigns storage locations, and performs auxiliary functions necessary to produce an executable machine-language program.

## COMPATIBILITY

The IBM System/360 Model 44 Programming System Assembler Language is a selected subset of the language available in the IBM System/360 programming support systems designed for the Models $30,40,50,65$, and 75 -- specifically:

- System 360 Operating System (os/360)
- System/360 Disk Operating System (DOS/360)
- System/360 Tape Operating System (TOS/360)

Thus, source programs written in the Model 44 assembler language can be assembled by the appropriate assembler, OS/360, DOS/360, or TOS/360, provided that (1) any source statements involving subroutine linkages or supervisory functions are modified to the format specified for the applicable system, (2) there are no statements using instructions peculiar to the Model 44 or the Model 44 assembler program, and (3) all SETA variable symbols are defined using the LCLA or GBLA statements as specified in the appropriate language. Appendix $G$ describes more specifically the relationship between the Model 44 Programming System Assembler

Language and the other System/360 programming support system assembler languages.

## THE ASSEMBIER T.ANGUAGE

The basis of the assembler language is a collection of mnemonic symbols that represent:

1. System/360 machine-language operation codes.
2. Operations (auxiliary functions) to be performed by the assembler program.

The language is augmented by other symbols, supplied by the programmer, and used to represent storage addresses or data. Symbols are easier to remember and to code than are their machine-language equivalents. Use of symbols greatly reduces programming effort and error.

## Machine Operation Codes

The assembler language provides mnemonic machine-instruction operation codes for all machine instructions that can be processed by the Model 44 source programs and extended mnemonic operation codes for the conditional branch instruction. Appendix D lists the acceptable machine operation codes for Model 44 source programs.

## Assembler Operation Codes

The assembler language also contains mnemonic assembler-instruction operation codes, used to specify auxiliary functions to be performed by the assembler. These are instructions to the assembler program itself and, with a few exceptions, do not result in the generation of machinelanguage code by the assembler program.

## THE ASSEMBLER PROGRAM

The assembler program, also referred to as the "assembler," processes the source statements written in the assembler language.

## Basic Functions

Processing a source program involves the translation of source statements into machine language, the assignment of storage locations to instructions and other elements of the program, and the performance of the auxiliary assembler functions designated by the programmer. The output of the assembler program is the relocatable module, a machine-language translation of the source program. The assembler furnishes a printed listing of the source statements and object program statements and additional information useful to the programmer in analyzing his program, such as error indications. The object program is in the format required by the linkage editor component of the System/360 Model 44 Programming System.

## PROGRAMMER AIDS

The assembler provides auxiliary functions that assist the programmer in checking and documenting programs, in controlling address assignments, in segmenting a program, in data and symbol definition, and in controlling the assembler itself. Mnemonic operation codes for these functions are provided in the language.

Variety in Data Representation: Decimal, hexadecimal, or character representation of machine-language binary values may be employed by the programmer in writing source statements. The programmer selects the representation best suited to his purpose.

Base Reqister Address Calculation: As discussed in the publication IBM System 360: Principles of Operation, Form A22-6821, the System/360 addressing scheme requires the designation of a base register (containing a base address value) and a displacement value in specifying a storage location. The assembler assumes the clerical burden of calculating storage addresses in these terms for the symbolic addresses used by the programmer. The programmer retains control of base register usage and the values entered there.

Relocatability: The object programs produced by the assembler are in a format enabling relocation from the originally assigned storage area to any other suitable area.

Sectioning and Linking: The assembler language and program provide facilities for partitioning an assembly into one or more parts called control sections. Control sections may be added or deleted when linkage editing the object program. Because control sections do not have to be loaded contiguously in storage, a sectioned program may be loaded and executed even though a continuous block of storage, large enough to accommodate the entire program, may not be available.

The assembler allows symbols to be defined in one assembly and referred to in another, thus effecting a link between separately assembled programs. This permits reference to data and transfer of control between programs. A detailed discussion of program sectioning and linking is contained in Section 3.

Program Listings: A listing of the source program statements and the resulting object program statements may be produced by the assembler for each source program it assembles. The programmer can control the form and content of the listing to some degree.

An alphabetical listing of all the symbols used in the program, together with cross references to the statements that use each symbol, can also be produced.

Error Indications: As a source program is assembled, it is analyzed for actual or potential errors in the use of the assembler language. Detected errors are indicated in the program listing, as described in the publication IBM System/360 Model 44 Programming System: Guide to System Use, Form C28-6812.

## PROGRAMMING SYSTEM RELATIONSHIPS

The assembler is a component of the IBM System 360 Model 44 Programming System and, as such, functions under control of the programming system. The programming system provides the assembler with input/output and other services needed in assembling a source program. In a like manner, the object program produced by the assembler will normally operate under control of the programming system and depend on it for input/output and other services. In writing the source program, the programmer uses the Supervisor call (SVC) instruction to invoke the facilities of the programming system supervisor. The programming system supervisor is discussed in the publication IBM System/360 Model 44 Programming System: Concepts and Facilities. Form C28-6810.

This section presents information about assembler language coding conventions, assembler source statement structure, and addressing.

## ASSEMBLER LANGUAGE CODING CONVENTIONS

This subsection discusses the general coding conventions associated with use of the assembler language.

## Coding Form

A source program is a sequence of source statements that are punched into cards. These statements may be written on the standard coding form, X28-6509 (shown in Figure 1), provided by IBM. One line of
coding on the form is punched into one card. The vertical columns on the form correspond to card columns. Space is provided on the form for program identification and instructions to keypunch operators. None of this information is punched into a card.

The body of the form (Figure 1) is composed of two fields: first, the statement field, columns 1-7i, and then the identification-sequence field, columns 73-80. The identification-sequence field is not part of a statement and is discussed following "Summary of Statement Format" in this section.

The entries (i.e., coding), composing a statement, occupy columns 1-71 of a line. Therefore, column 1 is referred to as the "begin" column and column 71 is referred to as the "end" column. The "begin" statement boundary may be altered by use of the Input Format Control (ICTL) assembler instruction


Figure 1. Coding Form
(discussed later) to designate an alternate begin column.

## Statement Format

A statement can be either a comment or an instruction.

A statement may be used for a comment by placing an asterisk in the begin column. Extensive comments entries may be written by using a series of lines with an asterisk in the begin column of each line.

Instructions may consist of one to four entries in the statement field. They are, from left to right: a name entry, an operation entry, an operand entry, and a comments entry. These entries must be separated by one or more blanks, and must be written in the order stated.

Only one statement is allowed per line; a statement cannot be continued on additional lines: column 72 must be blank. Columns 73 through 80 may contain a serial number, as discussed in Section 7, "Update Feature."

The coding form (Figure 1) is vertically ruled to provide an 8-character name field, a 5-character operation field, and a 56-character operand and/or comments field.

If desired, the programmer may disregard these column boundaries and write the name, operation, operand, and comments entries in other positions, subject to the following rules:

1. The entries must not extend beyond the "begin" and "end" statement boundaries (either the conventional boundaries, or the altered boundaries).
2. The entries must be in proper sequence, as stated above.
3. The entries must be separated by one or more blanks.
4. If used, a name entry must start in the begin column.

A description of the name, operation, operand, and comments entries follows:

Name Entries: The name entry is a symbol created by the programmer to identify a statement. A name entry usually is optional. The symbol must consist of eight characters or less, and be entered with the first character appearing in the begin column. If the begin column is blank, the assembler program assumes no name has been
entered. No blanks may appear in the name entry.

Operation Entries: The operation entry is the mnemonic operation code specifying the machine operation or assembler operation desired. An operation entry is mandatory and must start at least one position to the right of the begin column. Valid mnemonic operation codes for machine and assembler operations are contained in Appendixes D and $E$ of this publication. Valid operation codes consist of five characters or fewer for machine or assembler-instruction operation codes. No blanks may appear within the operation entry.

Operand Entries: The operand entry is the coding that identifies and describes data to be acted upon by the instruction, by indicating such things as storage locations, masks, storage-area lengths, or types of data.

Depending on the particular instruction, an operand entry may consist of one or more operands. Operands are required for all machine instructions but not for all assembler instructions.

Operands must be separated by commas, and no blanks may intervene between operands and the commas that separate them.

The operands may not contain embedded blanks, except as follows:

> If character representation is used to specify a constant, a literal, or immediate data in an operand, the character string may contain embedded blanks, e.g.. C'A D'.

Comments Entries: Comments are descriptive items of information about the program that are to be inserted in the program listing. All 256 valid characters (see "Character Set" in this section), including blanks, may be used in writing a comment. The entry must be separated from the operand entry by a blank. The comments entry cannot extend beyond column 71.

In statements where an optional operand entry is omitted but a comments entry is desired, the absence of the operand entry must be indicated by a comma preceded and followed by one or more blanks, as follows:


Statement Example: The following example illustrates the use of name, operation, operand, and comments entries. A compare instruction has been named by the symbol COMP; the operation entry (CR) is the mnemonic operation code for a register-toregister compare operation, and the two operands $(5,6)$ designate the two general registers whose contents are to be compared. The comments entry reminds the programmer that he is comparing "new sum" to "old" with this instruction.


## Summary of Statement Format

The entries in an instruction must always be in the following order: name, operation, operand(s), comment.

Every instruction requires an operation entry. Comments entries are optional. Name entries are required for certain instructions and are optional in all other cases. Operand entries are required for all machine instructions and most assembler instructions.

The name and operation entries must not contain embedded blanks. Operands must not have blanks preceding or following the commas that separate them.

All entries must be contained within the designated statement boundaries.

## Identification-Sequence Field

The identification-sequence field of the coding form (columns 73-80) is used to enter program identification and/or statement sequence characters. The entry is optional. If the field, or a portion of it, is used for program identification, the identification is punched in the source cards and reproduced in the program listing.

## Character Set

Source statements are written using the following characters:

Letters A through Z, and \$, \#, a
Digits $\quad 0$ through 9
Special
Characters + - = . * ( ) ' / E blank
These characters are represented by the card-punch combinations and internal bit configurations listed in Appendix A. In addition, any of the 256 punch combinations may be designated anywhere that characters are used in comments and between paired single quotes.

## ASSEMBLER LANGUAGE STRUCTURE

The basic structure of the language can be stated as follows.

A source statement is composed of:

- A name entry (usually optional).
- An operation entry (required).
- An operand entry (usually required).
- Comments entry (optional).

A name entry is:

- A symbol.

An operation entry is:

- A mnemonic operation code representing a machine or assembler instruction.

An operand entry is:

- One or more operands, each composed of one or more expressions. An expression is composed of a term or an arithmetic combination of terms.

Operands of machine instructions generally represent such things as storage locations, general registers, immediate data, or constant values. Operands of assembler instructions provide the information needed by the assembler program to perform the designated operation.

Figure 2 depicts this structure. Terms shown in Figure 2 are classed as absolute or relocatable. Terms are absolute or relocatable, depending on the effect of program relocation upon them. Program relocation is the loading of the object program into storage locations other than those originally assigned by the assembler. A term is absolute if its value does not change upon relocation. A term is relocatable if its value changes upon relocation.

The following subsection "Terms and Expressions" discusses these items as outlined in Figure 2.

TERMS AND EXPRESSIONS

TERMS

A term is a character or combination of characters that represents a value. This value may be assigned by the assembler (symbols, location counter reference) or may be inherent in the term itself (self-defining term, literal).

An arithmetic combination of terms is reduced to a single value by the assembler.

The following material discusses each type of term and the rules for its use.

## Symbols

A symbol is a character or combination of characters used to represent addresses or arbitrary values.

Symbols, through their use as names and in operands, provide the programmer with an efficient way to name and refer to a program element. A symbol, created by the programmer for use as a name entry and/or an operand, must conform to these rules:

1. The symbol must not consist of more than eight characters.
2. The first character must be a letter. The other characters may be letters. digits, or a combination of the two.
3. A symbol may not contain special characters, including blanks.

The following are valid symbols:

| READER | LOOP 2 | AB4 |
| :--- | :--- | :--- |
| A23456 | N | \$A1 |
| X4F2 | S4 | \#56 |

The following symbols are invalid, for the reasons noted:

| 256B | (first character is not a <br> letter) |
| :--- | :--- |
| RECORDAREA2 | (morecharacters) <br> BCD*34(contains the <br> character *) |
| IN AREA special |  |
| (contains a blank) |  |

DEFINING SYMBOLS: The assembler assigns a value to each symbol appearing as a name entry in a source statement. The value assigned to a symbol naming a storage area, an instruction, a constant, or a control section is the address of the leftmost byte of the storage field containing the named item. Since the address of such an item may change upon program relocation, the symbol naming it is considered a relocatable term.

A symbol used as a name entry in the Equate Symbol (EQU) assembler instruction is assigned the value designated in the operand entry of the instruction. Since the operand entry may represent a relocatable value or an absolute (i.e., nonchanging) value, the symbol is considered a relocatable term or an absolute term, depending upon the value to which it is equated.

The value of a symbol may not be negative and may not exceed 224-1.

A symbol is said to be defined when it appears as the name of a source statement. (A special case of symbol definition involving external references is discussed in Section 3, under "Program Sectioning and Linking.")

A symbol may be defined only once in an assembly. That is, each symbol used as the name of a statement must be unique within that assembly.

PREVIOUSLY DEFINED SYMBOLS: Some instructions require that a symbol appearing in the operand entry be previously defined. This simply means that the symbol, before its use in an operand, must have appeared as a name entry in a prior statement.

## Self-Defining Terms

A self-defining term is one whose value is inherent in the term. It is not assigned a value by the assembler. For example, the decimal self-defining term 15 represents a value of 15.


Figure 2. Assembler Language Structure -- Nachine and Assembler Instructions

There are three types of self-defining terms: decimal, hexadecimal, and character. Use of these terms is spoken of as decimal, hexadecimal, or character representation of a machine-language binary value or bit configuration.

Self-defining terms are classed as absolute terms, since the values they represent do not change upon program relocation.

USING SELF-DEFINING TERMS: Self-defining terms are the means of specifying machine values or bit configurations without equating the values to symbols and using the symbols.

Self-defining terms may be used to specify such program elements as immediate data, masks, registers, addresses, and address increments. The type of term selected (decimal, hexadecimal, or character) will depend on what is being specified.

The use of a self-defining term is quite distinct from the use of data constants or literals. When a self-defining term is used in a machine-instruction statement, its value is assembled into the instruction. When a data constant is referred to or a literal is specified in the operand of an instruction, its address is assembled into the instruction. self-defining terms are always right-justified; truncation or padding with zeros, if necessary, occurs on the left.

Decimal Self-Defining Term: A decimal self-defining term is simply an unsigned decimal number written as a sequence of decimal digits. High-order zeros may be used (e.g., 009). Limitations on the value of the term depend on its use. For example, a decimal term that designates a general register should have a value between 0 and 15; one that represents an address should not exceed the size of storage. In any case, a decimal term may not consist of more than eight digits; to be exact, it may not exceed 16777215 (224-1). A decimal self-defining term is assembled as its binary equivalent. Some examples of decimal self-defining terms are: 8, 147, 409?, and 00021.

Hexadecimal Self-Defining Term: A hexadecimal self-defining term is an unsigned hexadecimal number (written as a sequence | of one to six hexadecimal digits) enclosed in single quotes and preceded by the letter X: X'C49'.

Each hexadecimal digit is assembled as its 4-bit binary equivalent. Thus, a hexadecimal term used to represent an 8-bit mask would include two hexadecimal digits.

The maximum value of a hexadecimal term is FFFFFF.

The hexadecimal digits and their bit patterns are as follows:

| $0-0000$ | $4-0100$ | $8-1000$ | C- 1100 |
| :--- | :--- | :--- | :--- | :--- |
| $1-0001$ | $5-0101$ | $9-1001$ | D- 1101 |
| $2-0010$ | $6-0110$ | A- 1010 | E- 1110 |
| $3-0011$ | $7-0111$ | B- 1011 | F- 1111 |

A table for converting from hexadecimal representation to decimal representation is provided in Appendix B.

Character Self-Lefining Term: A character self-defining term consists of one to three characters enclosed by single quotes and preceded by the letter C. All letters, decimal digits, and special characters may be used in a character term. In addition, any of the remainder of the 256 punch combinations may be designated in a character self-defining term. Examples of character self-defining terms are as follows:

```
C'/' C' C'(blank)
C'ABC' C'13'
```

Because of the use of single quotes and ampersands as syntactic characters (ampersands are used as syntactic characters in variable symbols, which are discussed in Section 6), the following rule must be observed when using these characters in a character term:

For each single quote or ampersand desired in a character self-defining term, two single quotes or ampersands must be written. For example, the character values to the left are specified as indicated to the right:

$$
\begin{aligned}
& \text { A'\# C'A''\#' } \\
& \text { BEB C'BEEB' } \\
& \text { •' C'''•' } \\
& \text { 'g' C'''gछ'' }
\end{aligned}
$$

Each character in the character sequence is assembled as its 8 -bit code equivalent (see Appendix A). The two single quotes or ampersands that must be used to represent a single quote or ampersand within the character sequence are assembled as one single quote or ampersand.

## Location Counter Reference

The location counter reference enables the programmer to refer to the current value of the location counter. The location counter is used to assign storage addresses to program statements. It is the
assembler's equivalent of the instruction counter in the computer. As each machine instruction or data area is assembled, the location counter is first adjusted to the proper boundary for the item, if adjustment is necessary, and after the statement has been processed, incremented by the length of the assembled item. Thus, after a statement has been processed, it points to the next available location. If the statement is named by a symbol, the value attribute of the symbol is the value of the location counter after boundary adjustment. but before addition of the length.

For each successively declared control section, the location counter assigns locations in consecutively higher areas of storage. The first location of each control section is aligned to a double-word boundary. (Control sections are discussed further in Section 3, "Program Sectioning and Linking.")

The location counter setting can be controlled by using the START and ORG assembler instructions, which are described in Sections 3 and 5. The maximum value for the location counter is $2^{24-1}$.

The programmer may refer to the current value of the location counter at any place in a program by using an asterisk as a term in an operand. The asterisk represents the location of the first byte of currently available storage (i.e., after any required boundary adjustment). Using an asterisk as the operand in a machine-instruction statement is the same as placing a symbol in the name field of the statement and then using that symbol as an operand of the statement.

A reference to the location counter may be made in an address constant literal (i.e., the asterisk may be used in an address constant specified in literal form). The address of the instruction containing the literal is used for the value of the location counter. A location counter reference may not be used in a statement that requires the use of a predefined symbol, with the exception of the EQU and ORG assembler instructions.

## Literals

A literal may be used to introduce data into a program. It is simply a DC operand preceded by an equal sign (=).

A literal represents data rather than a reference to data. The appearance of a literal in a statement causes the assembler program to assemble the data specified by the literal, store this data in a "literal
pool," and place the address of the storage field containing the data in the operand field of the assembled statement.

Literals provide a means of entering constants (such as numbers for calculation, addresses, indexing factors, or words or phrases for printing out a message) into a program by specifying the constant in the operand of the instruction in which it is used. Specifying a literal is in contrast to using the DC assembler instruction to enter the data into the program and then specifying the name of the $D C$ instruction in the operand. only one literal is allowed in a machine-instruction statement .

A literal may not be combined with any other terms.

A literal may not be used as the receiving field of an instruction that modifies storage.

A literal may not be specified in an address constant (see Section 5, "DC--Define Constant").

The instruction coded below shows one use of a literal.


The statement GAMMA is a load instruction using a literal as the second operand. When assembled, the second operand of the instruction will be the address at which the value $\mathrm{F}^{\prime} 274^{\prime}$ is stored.

A literal may be used as an operand wherever a storage address is specified in a machine instruction or in a CCW assembler instruction operand. Literals are considered relocatable because the address of the literal, rather than the literal itself, will be assembled in the statement that employs a literal. The assembler generates the literals, collects them, and places them in a specific area of storage, as explained in the subsection "The Literal Pool." A literal is not to be confused with the immediate data in an SI instruction. Immediate data is assembled into the instruction.

Literal Format: The assembler requires a description of the type of literal being specified as well as the literal data itself. The descriptive portion of the literal must indicate the format of the constant.

The method of describing and specifying a constant as a literal is nearly identical to the method of specifying it in the operand of a DC assembler instruction. The major difference is that the literal must start with an equal sign (=), which indicates to the assembler that a literal follows. The reader is referred to the discussion of the DC assembler instruction operand format (section 5) for the means of specifying a literal. The type of literal designated in an instruction is not checked for correspondence with the operation code of the instruction.

Some examples of literals are:
\(\left.\begin{array}{lll}=A(B E T A) \& -- \& address constant literal <br>

=F^{\prime} 1234^{\prime} \& -- \& a fixed-point number with\end{array}\right]\)| a length of four bytes |
| :--- | :--- |

The Literal Pool: The literals processed by the assembler are collected and placed in a special area called the literal pool. and the location of the literal, rather than the literal itself, is assembled in the statement employing a literal. The positioning of the literal pool must be assigned by the programmer within the control section in which the literal is used.

The programmer may also specify that multiple literal pools be created. However, the sequence in which literals are ordered within the pool is controlled by the assembler. Further information on positioning the literal pool(s) is in Section 5 under "LTORG--Begin Literal Pool."

## EXPRESSIONS

This subsection discusses the expressions used in coding operand entries for source statements. Two types of expressions, absolute and relocatable, are presented along with the rules for determining these attributes of an expression.

As shown in Figure 2, an expression is composed of a single term or an arithmetic combination of terms. The arithmetic operators that may be used to combine the terms of an expression are + (addition), (subtraction), * (multiplication), and / (division).

The following are examples of valid expressions (provided that BETA, LAMBDA, GAMMA, TEN, and TWO are absolute):

| AREA $1+X^{\prime} 2 D^{\prime}$ | BETA*10 |
| :--- | :--- |
| $*+32$ | C'ABC' $^{\prime}$ |
| N-25 | 29 |
| FIELD | LAMBDA+GAMMA |
| $=F^{\prime} 1234^{\prime}$ | TEN/TWO |

The rules for coding expressions are:

1. An expression may not start with an arithmetic operator. Therefore, the expression -A+BETA is invalid. However, the expression $0-A+B E T A$ is valid.
2. An expression may not contain two terms or two operators in succession.
3. An expression may not consist of more than three terms.
4. An expression may not have more than one level of parentheses (i.e., a parenthetical expression may not appear within a parenthetical expression).
5. A multiterm expression may not contain a literal.
6. A parenthesized expression may not contain a literal.

## Evaluation of Expressions

A single term expression, e.g., 29, BETA, *, takes on the value of the term involved.

A multiterm expression, e.g., BETA+10, ENTRY-EXIT, $25 * 10+$ A, is reduced to a single value, as follows:

1. Each term is given its value.
2. Every expression is computed to 32 bits.
3. Arithmetic operations are performed left to right. Multiplication and division are done before addition and subtraction, e.g., A+B*C is evaluated as $A+(B * C)$, not $(A+B) * C$. The computed result is the value of the expression.
4. Division always yields an integer result; any fractional portion of the result is dropped. For example, 1/2*10 yields a zero result, whereas $10 * 1 / 2$ yields 5.
5. Division by zero is valid and yields a zero result.

A parenthesized multiterm expression used in an expression is processed before the rest of the terms in the expression, e.g., in the expression BETA* $(\operatorname{CON}-10)$, the
term CON-10 is evaluated first and the resulting value is used in computing the final value of the expression.

Negative values are carried in two's complement form. Final values of expressions are the truncated rightmost 24 bits of the results. The value of an expression before truncation must be in the range -224 through 224-1. A negative result is considered to be a 3-byte positive value. Intermediate results have a range of -23i through 231-1.

## Absolute and Relocatable Expressions

An expression is called absolute if its value is unaffected by program relocation.

An expression is called relocatable if its value changes upon program relocation.

The two types of expressions, absolute and relocatable, take on these characteristics from the term(s) composing them.

ABSOLUTE EXPRESSION: An absolute expression may be an absolute term or any arithmetic combination of absolute terms. An absolute term may be an absolute symbol, or any of the self-defining terms. All arithmetic operations are permitted between absolute terms.

An absolute expression may contain relocatable terms (RT) -- alone or in combination with absolute terms (AT) -- under the following conditions:

1. There must be an even number of relocatable terms in the expression.
2. The relocatable terms must be paired. Each pair of terms must have the same relocatability attribute, i.e., they appear in the same control section in this assembly (see section 3. "Program Sectioning and Linking"). Each pair must consist of terms with opposite signs. The paired terms do not have to be contiguous, e.g., RT+AT-RT.
3. No relocatable term may enter into a multiply or divide operation. Thus, RT-RT*10 is invalid. However, (RT-RT)*10 is valid.

The pairing of relocatable terms (with opposite signs and the same relocatability attribute) cancels the effect of relocation. Therefore, the value represented by the paired terms remains constant, regardless of program relocation. For example, in the absolute expression $A-Y+X, A$ is an absolute term, and $X$ and $Y$ are relocatable
terms with the same relocatability attribute. If A equals 50, Y equals 25, and $X$ equals 10 , the value of the expression would be 35. If $X$ and $Y$ are relocated by a factor of 100 , their values would then be 125 and 110. However, the value of the expression would still be 35 (50-125+110=35). An absolute expression reduces to a single absolute value.

The following examples illustrate absolute expressions. A is an absolute term; $X$ and $Y$ are relocatable terms with the same relocatability attribute.

```
A-Y+X
A
A*A
X-Y+A
```

*-Y (a reference to the location counter must be paired with another relocatable term from the same control section, i.e., with the same relocatability attribute)

RELOCATABLE EXPRESSION: A relocatable expression is one whose value would change by $n$ if the program in which it appears is relocated $\underline{n}$ bytes away from its originally assigned arēa of storage. All relocatable expressions must have a positive value.

A relocatable expression may be a relocatable term. A relocatable expression may contain relocatable terms -- alone or in combination with absolute terms -- under the following conditions:

1. There must be an odd number of relocatable terms.
2. All relocatable terms but one must be paired. Pairing is described in the preceding discussion of absolute expressions.
3. The unpaired term must not be directly preceded by a minus sign.
4. No relocatable term may enter into a multiply or divide operation.

A relocatable expression reduces to a single relocatable value. This value is the value of the odd relocatable term, adjusted by the values represented by the absolute terms and/or paired relocatable terms associated with it. The relocatability attribute is that of the odd relocatable term.

For example, in the expression $W-X+W$, the terms $W$ and $X$ are relocatable terms with the same relocatability attribute. If, initially, $W$ equals 10 and $x$ equals 15 , the value of the expression is 5. However, upon relocation, this value will change. If a relocation factor of 100 is applied,
the value of the expression is 105. Note that the value of the paired terms, $W-X$, remains constant at -5 regardless of relocation. Thus, the new value of the expression, 105, is the result of the value of the odd term (W) adjusted by the values of W -X.

The following examples illustrate relocatable expressions. A is an absolute
term, $W$ and $x$ are relocatable terms with the same relocatability attribute, $Y$ is a relocatable term with a different relocatability attribute.

| $\mathrm{Y}-32 * \mathrm{~A}$ | $=\mathrm{F}^{\prime} 1234^{\circ}$ (1iteral) |
| :--- | :--- |
| $\mathrm{W}-\mathrm{X}+\mathrm{Y}$ | $\mathrm{A} * \mathrm{~A}+\mathrm{W}$ |
| $\mathrm{W}-\mathrm{X}+*$ | $\mathrm{~W}-\mathrm{X}+\mathrm{W}$ |
| * (reference to | Y |
| $\quad$ location counter) |  |

## ADDRESSING

The IBM System/360 addressing technique requires the use of a base register, which contains the base address, and a displacement, which is added to the contents of the base register. The programmer may specify a symbolic address and request the assembler to determine its storage address composed of a base register and a displacement. The programmer may rely on the assembler to perform this service for him by indicating which general registers are available for assignment and what values the assembler may assume each contains. The programmer may use as many or as few registers for this purpose as he desires. The only requirement is that, at the point of reference, a register containing an address from the same control section is available, and that this address is less than or equal to the address of the item to which the reference is being made. The difference between the two addresses may not exceed 4095 bytes.

ADDRESSES -- EXPLICIT AND IMPLIED

An address is composed of a displacement plus the contents of a base register. (In the case of RX instructions, the contents of an index register are also used to derive the address in the machine.)

The programmer writes an explicit address by specifying the displacement and the base register number. In designating explicit addresses, a base register may not be combined with a relocatable symbol.

The programmer writes an implied address by specifying an absolute or relocatable address. The assembler has the facility to select a base register and compute a displacement, thereby generating an explicit address from an implied address, provided that it has been informed as to (1) what base registers are available to it and (2) what each contains. The programmer conveys this information to the assembler through the USING and DROP assembler instructions.

## BASE REGISTER INSTRUCTIONS

The USING and DROP assembler instructions enable programmers to use expressions representing implied addresses as operands of machine-instruction statements, leaving the assignment of base registers and the calculation of displacements to the assembler.

In order to use symbols in the operand field of machine-instruction statements, the programmer must (1) indicate to the assembler, by means of USING statements, which general registers are available for use as base registers, (2) specify, by means of the USING statement, what value each base register contains, and (3) load each base register with the value he has specified for it.

A program must have at least one USING statement for each control section that contains implicit addressing.

Having the assembler determine base registers and displacements relieves the programmer of separating each address into a displacement value and a base address value. This feature of the assembler will eliminate a likely source of programming errors, thus reducing the time required to check out programs. To take advantage of this feature, the programmer uses the USING and DROP instructions described in this subsection. The principal discussion of this feature follows the description of both instructions.

## USING -- Use Base Address Register

The USING instruction specifies a general register that is available for use as a base register. This instruction also states the base address value that the assembler may assume will be in the register at object time. Note that a USING instruction does not load the register specified. It is the programmer's responsibility to make sure that the specified base address value is placed into the register. Suggested loading methods are described in the subsection "Programming with the USING Instruction."

The format of the USING instruction statement is:


Operand $\underline{v}$ must be an absolute or relocatable expression. Literals are not permitted. Operand $v$ specifies a value that the assembler can use as a base address. The operand $r$ must be an absolute term. It specifies the general register that the assembler assumes will contain the base address represented by operand $\underline{v}$. The value of $\underline{r}$ must be in the range from 0 to 15.

For example, the following USING statement tells the assembler it may assume that the current value of the location counter will be in general register 12 at execution time.


If the programmer changes the value in a base register currently being used, and wishes the assembler to compute displacement from this value, the assembler must be told the new value by means of another USING statement. In the following sequence the assembler first assumes that the value of ALPHA is in register 9. The second statement then causes the assembler to assume that ALPHA+1000 is the value in register 9.


A USING statement may specify general register 0 as a base register if operand $v$ is a relocatable expression from any control section in the program or has an absolute value of zero. If general register 0 is specified, the assembler assumes that register 0 contains the value zero.

Note: If register 0 is made available by a USING instruction, the program is not relocatable, despite the fact that the value specified by operand $v$ must be relocatable. However, the programmer is able to make the program relocatable at some future time by:

1. Replacing register 0 with an alternate register in the USING statement.
2. Inserting an instruction that loads the alternate register with a relocatable value.
3. Reassembling the program.

## DROP -- Drop Base Register

The DROP instruction specifies a previously available register that may no longer be used as a base register. The format of the DROP instruction statement is as follows:


The absolute term indicates a general register previously named in a USING statement that is now unavailable for base addressing. The following statement, for example, prevents the assembler from using register 7:


It is not necessary to use a DROP statement when the base address being used is changed by a USING statement; nor are DROP statements needed at the end of the source program.

A register made unavailable by a DROP instruction can be made available again by a subsequent USING instruction.

## PROGRAMMING WITH THE USING INSTRUCTION

The USING (and DROP) instructions may be used anywhere in a program, as often as needed, to indicate the general registers that are available for use as base registers and the base address values that the assembler may assume each contains at exe-
cution time. Whenever an address is specified in a machine-instruction statement, the assembler determines whether there is an available register containing a suitable base address. A register is considered available for a relocatable address if it was specified in a USING instruction to have a relocatable value. A register with an absolute value is available only for absolute addresses. In either case, the base address is considered suitable only if it is less than or equal to the address of the item to which the reference is made. The difference between the two addresses may not exceed 4095 bytes. In calculating the base register to be used, the assembler will always use the available register giving the smallest displacement. If there are two registers with the same value, the highest numbered register will be used.


In the preceding sequence, the BALR instruction loads register 2 with the address of the immediately following storage location. In this case, it is the address of the instruction named FIRST. The USING instruction indicates to the assembler that register 2 contains this location. When employing this method, the USING instruction must immediately follow the BALR instruction. No other USING or load instructions are required if the location named LAST is within 4095 bytes of FIRST.

In the following sequence, the BALR and L instructions load registers 2 through 5 . The USING instructions indicate to the assembler that these registers are available as base registers for addressing a maximum of 16,384 consecutive bytes of storage, beginning with the location named HERE. The number of addressable bytes may be increased or decreased by changing the number of registers designated by the USING and $L$ instructions and the number of address constants specified in the DC instruction.

| \| Name | Operation | Operand |
| :---: | :---: | :---: |
| \| BEGIN | \| BALR | 2,0 |
|  | \| USING | \| HERE, 2 |
|  | \| USING | \| HERE+4096,3 |
|  | \| USING | \| HERE+8192,4 |
|  | \| USING | \| HERE+12288,5 |
| \| HERE | \| L | 3, BASEAD |
|  | \| L | 14, BASEAD+4 |
|  | \| L | 5, BASEAD+8 |
|  | \| $\mathrm{B}^{\text {}}$ | FIRST |
| BASEAD | \| DC | A (HERE+4096) |
|  | \| DC | A ( $\mathrm{HERE}+8192$ ) |
|  | \| DC | \| A (HERE+12288) |
| \|FIRST | i . |  |
|  | \| . |  |
|  | 1. |  |
| \|LAST | \| . |  |
|  | \| END | BEGIN |

## RELATIVE ADDRESSING

Relative addressing is the technique of addressing instructions and data areas by designating their location in relation to the location counter or to some symbolic location. This type of addressing is always in bytes, never in bits, words, or instructions. Thus, the expression *+4 specifies an address that is four bytes greater than the current value of the location counter. In the sequence of instructions shown in the following example, the location of the $C R$ machine instruction can be expressed in two ways, ALPHA 2 or BETA-4, because all of the mnemonics in the example are for 2-byte instructions in the RR format.


## PROGRAM SECTIONING AND LINKING

It is often convenient, or necessary, to write a program in sections. The sections may be assembled separately, then combined
via the linkage editor into one or more executable phases. The assembler provides facilities for creating multisectioned programs and symbolically linking separately assembled programs or program sections. The combined number of control sections and dummy sections plus the number of unique symbols in EXTRN statements may not exceed 255.

Sectioning a program is optional, and many programs can best be written without sectioning. The programmer writing an unsectioned program need not concern himself with the subsequent discussion of program sections, which are called control sections. He need not employ the CSECT instruction, which is used to identify the control sections of a multisection program. Similarly, he need not concern himself with the discussion of symbolic linkages if his program neither requires a linkage to nor receives a linkage from another program. He may, however, wish to identify the program and/or specify a tentative starting location for it, both of which may be done by using the START instruction. He may also want to employ the dummy section feature obtained by using the DSECT instruction.

Note: Program sectioning and linking is closely related to the specification of base registers for each control section. Sectioning and linking examples are provided under the heading "Addressing External Control Sections."

## CONTROL SECTIONS

The concept of program sectioning should be taken into consideration at coding time, assembly time, and load time. To the programmer, a program is a logical unit. He may want to divide it into sections called control sections; if so, he writes it in such a way that control passes properly from one section to another regardless of the relative physical position of the sections in storage. A control section is a block of coding that can be relocated independently (i.e., without affecting the location of other coding), at load time, without altering or impairing the operating logic of the program. It is normally identified by the CSECT instruction. However, if it is desired to specify a tentative starting location, the START instruction may be used to identify the first control section.

To the assembler, there is no such thing as a program; instead, there is an assembly, which consists of one or more control sections. (However, the terms assembly and
program are often used interchangeably.) An unsectioned program is treated as a single control section. To the linkage editor, there are no programs, only control sections that must be fashioned into an object program.

The assembler output consists of the assembled control sections and a control dictionary. The control dictionary contains information the linkage editor needs to complete cross-referencing between control sections as they are combined into an object program. The linkage editor can combine control sections from various assemblies with the help of the corresponding control dictionaries. Successful combination of separately assembled control sections depends on the techniques used to provide symbolic linkages between the control sections. Whether the programmer writes an unsectioned program, a multisectioned program, or part of a multisectioned program, he still knows what eventually will be entered into storage because he has described storage symbolically. He may not know where each section appears in storage, but he does know what storage contains. There is no constant relationship between control sections. Thus, knowing the location of one control section does not make another control section addressable by relative addressing techniques.

## Control Section Location Assignment

Locations are assigned to control sections as if the sections are placed in storage consecutively, in the same order as they first occur in the program. Each control section subsequent to the first begins at the next available double-word boundary.

## START -- Start Assembly

The START instruction may be used to give a name to the first (or only) control section of a program. It may also be used to specify an initial location counter value for the program. The format of the START instruction statement is as follows:


If a symbol names the START instruction, the symbol is established as the name of the control section. Otherwise, the control section is considered to be unnamed. All subsequent statements are assembled as part of that control section. This continues until an instruction identifying a different control section (CSECT, DSECT, or COM) is encountered.

The symbol in the name field is a valid relocatable symbol whose value represents the address of the first byte of the control section.

The assembler uses the self-defining term specified by the operand as the initial location counter value of the program. This value should be divisible by eight. For example, either of the following statements could be used to assign the name PROG2 to the first control section and to indicate an initial assembly location of 2040 .


If the operand is omitted, the assembler sets the initial location counter value of the program at zero. The location counter is set at the next double-word boundary when the value of the START operand is not divisible by eight.

Note: The START instruction may not be preceded by any type of assembler language statement that may either affect or depend upon the setting of the location counter.

## CSECT -- Identify control Section

The CSECT instruction identifies the beginning of a control section. The format of the CSECT instruction statement is as follows:


If a symbol names the CSECT instruction, the symbol is established as the name of the control section; otherwise, the section is considered to be unnamed. Multiple CSECT instructions must have unique names.

The name of a CSECT may be blank, provided that no other CSECT or START instruction has a blank name.

The symbol in the name field is a valid relocatable symbol whose value represents the address of the first byte of the control section.

The occurrence of a CSECT instruction terminates the previous control section.

## Unnamed Control Section

If it is desired to write a program that is unsectioned, the program does not need to contain a CSECT or START instruction. In this case, the assembler will generate an unnamed START statement for the first assembler language statement that may either affect or depend upon the setting of the location counter.

## DSECT -- Identify Dummy Section

A dummy section represents a control section that is assembled but is not part of the object program. A dummy section is a convenient means of describing the layout of an area of storage without actually reserving the storage. (It is assumed that the storage is reserved either by some other part of this assembly or else by another assembly.) The DSECT instruction identifies the beginning of a dummy section. More than one dummy section may be defined per assembly, but each must be named. The format of the DSECT instruction statement is as follows:


The symbol in the name field is a valid relocatable symbol whose value represents the first byte of the section.

All statements following the DSECT instruction are assembled as part of that control section until a statement identifying a different control section is encountered (i.e., another DSECT, CSECT, or COM instruction). All assembler language instructions may occur within dummy sections.

Symbols that name statements in a dummy section may be used in USING instructions. Therefore, they may be used in program elements (e.g., machine-instructions and data definitions) that specify storage addresses. An example illustrating the use of a dummy section appears subsequently under "Addressing Dummy Sections."

The occurrence of a DSECT instruction terminates the previous control section. A DSECT cannot be resumed.

Note: A symbol that names a statement in a dummy section may be used in an A-type address constant only if it is paired with another symbol (with the opposite sign) from the same dummy section.

DUMMY SECTION LOCATION ASSIGNMENT: A location counter is used to determine the relative locations of named program elements in a dummy section. The location counter is always set to zero at the beginning of the dummy section, and the location values assigned to symbols that name statements in the dummy section are relative to the initial statement in the section.

ADDRESSING DUMMY SECTIONS: The programmer may wish to describe the format of an area whose storage location will not be determined until the program is executed. He can describe the format of the area in a dummy section, and he can use symbols defined in the dummy section as the operands of machine instructions. To effect references to the storage area, he does the following:

1. Provides a USING statement specifying both a general register that the assembler can assign to the machineinstructions as a base register and a value from the dummy section that the assembler may assume the register contains.
2. Ensures that the same register is loaded with the actual address of the storage area.

The values assigned to symbols defined in a dummy section are relative to the initial statement of the section. Thus, all machine instructions which refer to names defined in the dummy section will, at execution time, refer to storage locations relative to the address loaded into the register.

| Name | \|Operation | Operand |
| :---: | :---: | :---: |
| ASMBL2 | \|CSECT | \| |
| \| BEGIN | \| BALR | 12,0 |
|  | \| USING | 1*. 2 |
|  | 1. | 1 |
|  | - | , |
|  | 1. |  |
| , | \| USING | \|INAREA, 3 |
|  | \|CLI | \| INCODE, ${ }^{\prime}{ }^{\text {A }}$ |
|  | \|BE | \|ATYPE |
|  | - |  |
|  | 1. | 1 |
|  | 1. |  |
| \| ATYPE | \|LA | 15,0 |
|  | 1 LA | 16.5 |
| I | \|L | 17.INPUTA(5) |
| , | \|ST | 17.WORKA (5) |
| 1 | \| LA | 15.4(5) |
| , | \|BCT | $16, *-12$ |
| 1 | \| LA | 15.0 |
| I | \| LA | 16.9 |
| I | \| LH | 17.INPUTB(5) |
| I | \|STH | 17,WORKB(5) |
| 1 | \|LA | 15,2(5) |
| + | \| BCT | 16,*-12 |
| I | 1. |  |
| 1 | 1. | 1 |
|  | $1 \cdot$ |  |
| \| WORKA | \|DS | \| 5 F |
| \| WORKB | \| DS | 19 H |
| 1 | 1 - | , |
| 1 | 1. |  |
| I | \| . | , |
| \| INAREA | \| DSECT | \| |
| 1 INCODE | \|DS | \|CL1 |
| \| INPUTA | \|DS | \|5F |
| 1 INPUTB | \| DS | 19 H |
| 1 | 1. | 1 |
| 1 | 1. | 1 |
| 1 | 1 - | 1 |
| 1 | \| END | 1 |

An example of addressing dummy sections is shown in the foregoing coding. Assume that two independent assemblies (Assembly 1 and Ássembiy 2 ) have been loaded and are to be executed as a single overall program. Assembly 1 is an input routine that places a record in a specified area of storage, places the address of the input area containing the record in general register 3, and branches to Assembly 2. Assembly 2 processes the record. The coding shown in the example is from Assembly 2.

The input area is described in Assembly 2 by the DSECT control section INAREA. Portions of the input area (i.e., record) that the programmer wishes to work with are named in the DSECT control section as shown. The assembler instruction USING INAREA, 3 designates general register 3 as the base register to be used in addressing the DSECT control section, and it is assumed that general register 3 contains the address of INAREA.

Assembly 1, during execution, loads the actual beginning address of the input area in general register 3 . Because the symbols used in the DSECT section are defined relative to the initial statement in the section, the address values they represent will, at the time of program execution, be the actual storage locations of the input area.

## COM -- DEFINE COMMON CONTROL SECTION

The com assembier instruction identifies and reserves a common area of storage that may be referred to by independent assemblies that have been linked and loaded for execution as one overall program. The format is:


Only one common control section can be designated in an assembly.

When several assemblies, each designating a common control section of the same name, are linkage edited, the amount of storage reserved for this name is equal to the longest of these common control sections. (In this context, a blank common control section is considered to be uniquely named.)

The common area may be broken up into subfields through use of the DS and DC assembler instructions. Names of subfields are defined reiative to the beginning of the common section, as in the DSECT control section.

Instructions or constants that appear in a common control section are not assembled, i.e., no machine language code is generated for them. As much storage is reserved as would be required for the instructions or constants if they were assembled. Data can be placed in a common control section only through execution of the program.

If the assignment of common storage is done in the same manner by each independent assembly, reference to a location in common by any assembly results in the same location being referenced. When assembled, common location assignment starts at zero.

The occurrence of a COM instruction terminates the previous control section.


The above statements reserve a common area of storage that is 42 bytes long. The common area contains three subfields: AREA1 occupies five fullwords ( 20 bytes), AREA2 occupies nine halfwords (18 bytes), and MASK occupies 4 bytes. No machine language code is generated.

## SYMBOLIC LINKAGES

Symbols may be defined in one program and referred to in another, thus effecting symbolic linkages between independently assembled programs. The linkages can be effected only if the assembler is able to provide information about the linkage symbols to the linkage editor, which resolves these linkage references during a subsequent phase of processing. The assembler places the necessary information in the control dictionary on the basis of the linkage symbols identified by the ENTRY and EXTRN instructions. Note that these symbolic linkages are described as linkages between independent assemblies; more specifically, they are linkages between independently assembled control sections.

In the program where the linkage symbol is defined (i.e., used as a name), it must
also be identified to the assembler by means of the ENTRY assembler instruction. It is identified as a symbol that names an entry point, which means that another program may use that symbol in order to effect a branch operation or a data reference. The assembler places this information in the control dictionary.

Similarly, the program that uses a symbol defined in some other program must identify it by the EXTRN assembler instruction. It is identified as an externally defined symbol (i.e.. defined in another program) that is used to effect linkage to the point of definition. The assembler places this information in the control dictionary.

## ENTRY -- IDENTIFY ENTRY-POINT SYMBOL

The ENTRY instruction identifies linkage symbols that are defined in this program but may be used by some other program. The format of the entry instruction statement is as follows:


A program may contain a maximum of 100 ENTRY symbols. ENTRY symbols that are not defined (i.e., that do not appear as statement names), although invalid, will also count toward this maximum of 100 ENTRY symbols.

The symbol in the ENTRY operand field may be used as an operand by another program. An ENTRY statement operand may not contain a symbol defined in a dummy section or common. The following example identifies the statement named SINE as an entry point to the program.


Note: The name of a control section does not have to be identified by an ENTRY instruction when another program uses it as an entry point. The assembler automatically places information on control section names in the control dictionary.

## EXTRN -- IDENTIFY EXTERNAL SYMBOL

The EXTRN instruction identifies linkage symbols that are used by this program but defined in some other program. Each external symbol must be identified; this includes symbols that name control sections. The format of the EXTRN instruction statement is as follows:


The symbol in the operand field may not appear as a name of a statement in this program. It may not be used in expressions requiring that all symbols be previously defined. Thus, an EXTRN symbol may not be used in the operand of an EQU assemblex instruction.

The following example identifies two external symbols that have been used as operands in this program but are defined in some other program.


An example that employs the EXTRN instruction appears subsequently under "Addressing External Control Sections."

Note: When external symbols are used in an expression, they may not be paired. Each external symbol must be considered as having a unique relocatability attribute.

## Addressing External Control Sections

One way in which a program is linked to an external control section is to have the program:

1. Identify the external symbol with the EXTRN instruction and create an address constant from the symbol.
2. Load the address constant into a general register and branch to the control section via the register.


For example, to link to the control section named SINE, the preceding coding might be used.

An external symbol naming data may be referred to as follows:

1. Identify the external symbol with the EXTRN instruction, and create an address constant from the symbol.
2. Load the constant into a general register, and use the register for base addressing.

For example, to use an area named RATETB, which is in another control section, the following coding might be used:


The combined number of control sections and dummy sections plus the number of unique symbols in EXTRN statements may not exceed 255.

This section discusses the coding of the machine instructions represented in the assembler language. The reader is reminded that the functions of each machine instruction are discussed in the publication IBM System/360: Principles of operation, Form A22-6821.

This section should be used in conjunction with Appendix $C$, which describes assembler operand field formats for the various machine instructions.

## MACHINE INSTRUCTION STATEMENTS

Machine instructions may be represented symbolically as assembler language statements. The symbolic format of each varies according to the actual machine instruction format. Four formats are acceptable to the assembler: RR, RX, RS, and SI. Within each basic format, further variations are possible.

The symbolic format of a machine instruction is similiar to, but does not duplicate, its actual format. Appendix $C$ illustrates machine format for the four classes of instructions. A mnemonic operation code is written in the operation field, and one or more operands are written in the operand field. Comments may be appended to a machine instruction statement as previously explained in Section 1.

Any machine instruction statement may be named by a symbol, which other assembler statements can use as an operand. The value of the symbol is the address of the leftmost byte assigned to the assembled instruction.

## Instruction Alignment and Checking

All machine instructions are aligned automatically by the assembler on halfword boundaries. If any statement that causes information to be assembled requires alignment, the bytes skipped are filled with hexadecimal zeros. All expressions that specify storage addresses are checked to ensure that they refer to appropriate boundaries for the instructions in which they are used. Register numbers are also checked to make sure that they specify the proper registers, as follows:

1. Floating-point instructions must specify floating-point registers 0, 2, 4, or 6.
2. Double-shift, fullword multiply, and divide instructions must specify an even-numbered general register in the first operand.

## OPERAND FIELDS AND SUBFIELDS

Some symbolic operands are written as a single field, and other operands are written as a field followed by one or two subfields. For example, addresses consist of the contents of a base register and a displacement. An operand that specifies a base and displacement is written as a displacement field followed by a base register subfield, as follows: 40(5). In the RX format, an operand that specifies both an index register and a base register is written as follows: 40(3,5).

Appendix $C$ shows two types of addressing formats for RX, RS, and SI instructions. In each case, the first type shows the method of specifying an address explicitly, as a base register and displacement. The second type indicates how to specify an implied address as an expression.

For example, an Add instruction (RX format) may have either of the following symbolic operands:

R1,D2(X2,B2) -- explicit address
R1,S2 (X2) -- implied address
Whereas D2 must be represented by an absolute expression, $S 2$ may be represented by either a relocatable or an absolute expression. Both X 2 and B 2 must be absolute terms.

In order to use implied addresses, the following rules must be observed:

1. The base register assembler instructions (USING and DROP) must be used.
2. An explicit base register designation must not accompany the implied address.

For example, assume that FIELD is a relocatable symbol that has been assigned a value of 7400. Assume also that the assembler has been notified (by a USING
instruction) that general register 12 currently contains a relocatable value of 4096 and is available as a base register. The following example shows a machine instruction statement as it would be written in assembler language and as it would be assembled. Note that the value of D2 is the difference between 7400 and 4096 and that $X 2$ is assembled as zero, since it was omitted. The assembled instruction is presented in hexadecimal:

Assembler statement:

$$
\text { ST } \quad 4, \text { FIELD }
$$

Assembled instruction:

$$
\begin{array}{llllr}
\text { Op. Code } & \text { R1 } & \text { X2 } & \text { B2 } & \text { D2 } \\
50 & 4 & 0 & \text { C } & \text { CE8 }
\end{array}
$$

An address may be specified explicitly as a base register and displacement (and index register for $R X$ instructions) by the formats shown in the second column of Table 1. The address may be specified as an implied address by the formats shown in the third column.

Table 1. Details of Address Specification

| \| Type | \|Explicit Address | \| Implied Address |
| :---: | :---: | :---: |
| $\mathrm{j}_{\mathrm{RX}}$ | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | 1S2(x2) |
|  | D2 (, B2) | \|S2 |
| \|RS | D2 (B2) | \|S2 |
| \|SI | D1 (B1) | \|S1 |

A comma must separate operands. Parentheses must enclose a subfield or subfields, and a comma must separate two subfields within parentheses. When parentheses are used to enclose one subfield, and the subfield is omitted, the parentheses must be omitted. In the case of two subfields that are separated by a comma and enclosed by parentheses, the following rules apply:

1. If both subfields are omitted, the separating comma and the parentheses must also be omitted.
```
L 2,48(4,5)
L 2,FIELD
```

(no indexing, implied address)
2. If the first subfield in the sequence is omitted, the comma that separates it from the second subfield is written. The parentheses must also be written.

L $\quad 2.48(4,5)$
L 2,48(,5) (no indexing)
3. If the second subfield in the sequence is omitted, the comma that separates it from the first subfield must be omitted. The parentheses must be written.

$$
\mathrm{L} \quad 2,48(4,5)
$$

L 2,FIELD (4)
(implied address)
Fields and subfields in a symbolic operand may be represented either by absolute or by relocatable expressions, depending on what the field requires. (An expression has been defined as consisting of one term or a series of arithmetically combined terms.) Refer to Appendix $C$ for a detailed description of field requirements.

Note: Blanks may not appear in an operand unless provided by a character selfdefining term or a character literal. Thus, blanks may not intervene between fields and the comma separators, between parentheses and fields, etc.

## MACHINE-INSTRUCTION MNEMONIC CODES

The mnemonic operation codes (shown in Appendix D) are designed to be easily remembered codes that indicate the functions of the instructions. The normal format of the code is shown below; the items in brackets are not necessarily present in all codes:

## Verb[Modifier] [Data Type] [Machine Format]

The verb, which is usually one or two characters, specifies the function. For example, A represents Add, and $S T$ represents Store. The function may be further defined by a modifier. For example, the modifier $L$ indicates a logical function, as in AL for Add Logical.

Mnemonic codes for functions involving data usually indicate the data types by letters that correspond to those for the data types in the DC assembler instruction (see Section 5). Furthermore, letters U and $W$ have been added to indicate, respectively, short and long, unnormalized floating-point operations, and letters D and $E$ have been added to indicate, respectively, long and short, normalized floating-point operations. For example, AE indicates Add Normalized Short, whereas AW indicates Add Unnormalized Long. Where applicable, fullword fixed-point data is implied if the data type is omitted.

The letters $R$ and $I$ are added to the codes to indicate, respectively, RR and SI machine instruction formats. Thus, AER indicates Add Normalized Short in the RR format.

The examples that follow are grouped according to machine instruction format. They illustrate the various symbolic operand formats. All symbols employed in the examples must be assumed to be defined elsewhere in the same assembly. All symbols that specify register numbers and lengths must be assumed to be equated elsewhere to absolute values.

Implied addressing, control section addressing, and the function of the USING assemoler instruction are not considered here. For discussion of these considerations and for examples of coding sequences that illustrate them, the reader is referred to Section 3, "Program sectioning and Linking" and "Base Register Instructions."

## RR Format

| \| Name | Operation | Operand |
| :---: | :---: | :---: |
| \| ALPHA1 | \| LR | 11,2 |
| \| ALPha2 | \| LR | \| REG1, REG2 |
| \| BETA | SPM | 115 |
| \|GAMMA1 | ISVC | 1250 |
| \| GAMMA2 | \|SVC | \| TEN |

The operands of ALPHA1, BETA, and GAMMA1 are decimal self-defining values that are categorized as absolute expressions. The operands of ALPHA2 and GAMMA2 are symbols that are equated elsewhere to absolute values.

## RS Format

| \| Name | \|Operation | Operand |
| :---: | :---: | :---: |
| \| ALPHA1 | \|SLL | \| REG2, 15 |
| \| ALPHA2 | \|SLL | \| REG2,0(15) |

ALPHA1 is a shift instruction shifting the contents of REG2 left 15 bit positions. ALPHA2 is a shift instruction shifting the contents of REG2 left by the value contained in general register 15 .

RX Format

| \| Name | \|Operation | Operand |
| :---: | :---: | :---: |
| \| ALPHA1 | 1 L | [1,39(4,10) |
| \| ALPHA2 | \| L | \|REG1,39(4,TEN) |
| \| BETA1 | \| L | \|2, ZETA (4) |
| \| BETA2 | \| L | \| REG2, 2ETA (REG4) |
| \| GAMMA1 | \| L | \| 2, ZETA |
| \| GAMMA2 | \| | \| REG2, ZETA |
| \| GAMMA3 | \| L | \| $2,=\mathrm{F}^{\prime} 1000^{\prime}$ |
| \| LAMBDA | \| | 13,20(,5) |

Both ALPHA instructions specify explicit addresses; REG1 and TEN are absolute symbols. Both BETA instructions specify implied addresses, and both use index registers. Indexing is omitted from the GAMMA instructions. GAMMA1 and GAMMA2 specify implied addresses. The second operand of GAMMA3 is a literal. LAMBDA specifies no indexing.

SI Format

| Name | \|operation | \|Operand |
| :---: | :---: | :---: |
| \| ALPHA1 | \|CLI | \|40(9). X'40' |
| \| ALPHA2 | \|CLI | \| 40 (REG9) ,TEN |
| \|BETA1 | \|CLI | \| ZETA, TEN |
| \| BETA2 | \| CLI | \| ZETA, C'A' |
| GAMMA1 | \|SIO | 140(9) |
| \|GAMMA2 | ISIO | 10(9) |
| \|GAMMA3 | \|SIO | 140(0) |
| GAMMA4 | \|SIO | \| ZETA |

The ALPHA instructions and GAMMA1 through GAMMA3 instructions specify explicit addresses, whereas the BETA instructions and GAMMA4 instruction specify implied addresses. GAMMA2 specifies a displacement of zero. GAMMA3 specifies no base register.

## EXTENDED MNEMONIC CODES

For the convenience of the programmer, the assembler provides extended mnemonic codes, which allow conditional branches to be specified mnemonically as well as through the use of the BC machineinstruction. These extended mnemonic codes specify both the machine branch instruction and the condition on which the branch is to occur. The codes are not part of the universal set of machine instructions, but are translated by the assembler into the

| \| Exten | ded code | Meaning | Machine-Instruction |
| :---: | :---: | :---: | :---: |
| \|B | D2 ( $22, \mathrm{~B} 2)$ | Branch Unconditional | BC 15, D2 (X2, B2) |
| \| PR | R2 | Branch Unconditional (RR format) | BCR 15,R2 |
| \| NOP | D2 (X2, B2) | No Operation | BC 0,D2 (X2,B2) |
| \| NOPR | R2 | No Operation (RR format) | BCR 0,R2 |
| Used After Compare Instructions |  |  |  |
| \| BH | D2 (X2,B2) | Branch on High | BC 2,D2 (X2,B2) |
| \| BL | D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ | Branch on Low | BC 4,D2 $\mathrm{X} 2, \mathrm{~B} 2)$ |
| $\mid \mathrm{BE}$ | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Equal | BC 8,D2 ${ }^{\text {( } 2, B 2 \text { ) }}$ |
| \| BNH | D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ | Branch on Not High | BC 13,D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ |
| \| BNL | D2 (x2, B2) | Branch on Not Low | BC 11, D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ |
| \| BNE | D2 (X2, B2) | Branch on Not Equal | BC 7,D2 (X2,B2) |
| Used After Arithmetic Instructions |  |  |  |
| \| BO | D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ | Branch on Overflow | BC 1,D2 (X2,B2) |
| \|BP | D2 (X2, B2) | Branch on Plus | BC 2,D2 (X2,B2) |
| \| BM | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Minus | BC 4,D2 ${ }^{\text {( } 22, B 2 \text { ) }}$ |
| \| BZ | D2 (X2,B2) | Branch on zero | BC 8,D2 (X2,B2) |
| \| BNP | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Not Plus | BC 13, D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ |
| \| BNM | D2 (x2, B2) | Branch on Not Minus | BC 11,D2 (X2,B2) |
| \| BNZ | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Not Zero | BC 7,D2 ${ }^{\text {(X2,B2) }}$ |
| Used After Test Under Mask Instructions |  |  |  |
| \| BO | D2 (X2, B2) | Branch if Ones | BC 1, D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ |
| \| $\mathrm{BM}^{\text {m }}$ | D2 (X2, B2) | Branch if Mixed | BC 4,D2 $\mathrm{X} 2, \mathrm{~B} 2)$ |
| \| BZ | D2 (X2, B2) | Branch if zeros | BC 8,D2 (X2,B2) |
| \| BNO | D2 (X2, B2) | Branch if Not Ones | BC 14, D2 (X2,B2) |

Figure 3. Extended Mnemonic Codes
corresponding operation and condition combinations.

The allowable extended mnemonic codes and their operand formats are shown in Figure 3, together with their machineinstruction equivalents. Unless otherwise noted, all extended mnemonics shown are for instructions in the RX format. Note that the only difference between the operand fields of the extended mnemonics and those of their machine instruction equivalents is the absence of the R1 field and the comma that separates it from the rest of the operand field. The extended mnemonic list, like the machine instruction list, shows explicit address formats only. Each address can also be specified as an implied address.

In the following examples, which illustrate the use of extended mnemonics, it is to be assumed that the symbol GO is defined elsewhere in the program.


The first two instructions specify an unconditional branch to an explicit address. The address in the first case is the sum of the contents of base register 6 , the contents of index register 3 , and the displacement 40; the address in the second instruction is not indexed. The third instruction specifies a branch on low to the address implied by $G O$ as indexed by the contents of index register 3; the fourth instruction does not specify an index register. The last instruction is an unconditional branch to the address contained in register 4.

Just as machine instructions are used to request the computer to perform a sequence of operations during program execution time, so assembler instructions are requests to the assembler to perform certain operations during the assembly. Assembler-instruction statements, in contrast to machine-instruction statements, do not always cause machine instructions to be included in the assembled program. Some, such as DS and DC, generate no instructions but do cause storage areas to be set aside for constants and other data. Others, such as EQU and SPACE, are effective only at assembly time; they generate nothing in the assembled program and have no effect on the location counter.

The following is a list of assembler instructions.

> Symbol Definition Instruction EQU -- Equate Symbol Data Definition Instructions DC -- Define Constant DS -- Define Storage CCW -- Define Channel command word  * Program Sectioning and Linking Instructions START -- Start Assembly CSECT -- Identify Control Section DSECT -- Identify Dummy Section ENTRY -- Identify Entry-Point Symbol EXTRN -- Identify External Symbol COM -- Identify Common Control Section

* Base Register Instructions

USING -- Use Base Address Register
DROP -- Drop Base Address Register
Listing Control Instructions
TITLE - Identify Assembly Output
EJECT -- Start New Page
SPACE -- Space Listing
PRINT -- Print Optional Data
Program Control Instructions
ICTL -- Input Format Control
ORG -- Set Location Counter
LTORG -- Begin Literal Pool
CNOP -- Conditional No Operation
END -- End Assembly
REPRO -- Reproduce Following Card

[^1]SYMBOL DEFINITION INSTRUCTION

EQU -- EQUATE SYMBOL

The EQU instruction is used to define a symbol by assigning to it the value and the relocatability attribute of an expression in the operand field. The format of the EQU instruction statement is as follows:


The expression in the operand field may be absolute or relocatable. Any symbols appearing in the expression must be previously defined.

The symbol in the name field is given the value and the relocatability attribute of the expression in the operand field.

The EQU instruction is the means of equating symbols to register numbers, immediate data, and other arbitrary values. The following examples illustrate how this might be done:


To reduce programming time, the programmer can equate symbols to frequently used expressions and then use the symbols as operands in place of the expressions. For example:


The name, AREA, is defined as ALPHA-BETA+GAMMA and may be used in place of it. Note, however, that ALPHA, BETA, and GAMMA must all be previously defined.

## DATA_DEFINITION INSTRUCTIONS

There are three data definition instruction statements: Define constant (DC), Define Storage (DS), and Define Channel Command Word (CCW).

These statements are used to enter data constants into storage, to define and reserve areas of storage, or to specify the contents of channel command words. The statements may be named by symbols so that other program statements can refer to the fields generated from them.

DC -- DEFINE CONSTANT

The DC instruction is used to provide constant data in storage. A variety of constants may be specified: fixed-point, floating-point, hexadecimal, character, and storage addresses. (Data constants are generally called constants unless they are created from storage addresses, in which case they are called address constants.) The format of the DC instruction statement is as follows:


The operand consists of four subfields: the first three describe the constant; the fourth provides the constant or constants. The first and third subfields may be omitted, but the second and fourth must be specified. Note that more than one constant may be specified in the fourth subfield for most types of constants. Each constant so specified must be of the same type; the descriptive subfields that precede the constants apply to all of them.

No blanks may occur within any of the subfields (unless provided as characters in a character constant or a character selfdefining term), nor may they occur between the subfields of an operand.

The subfields of the DC operand are written in the following sequence:


The symbol that names the DC instruction is the name of the constant (or first constant if the instruction specifies more than one). Relative addressing ie.g., SYMBOL+2) may be used to address the various constants if more than one has been specified, because the number of bytes allocated to each constant can be determined.

The value of the symbol naming the DC instruction is the address of the leftmost byte (after alignment) of the first, or only, constant.

Boundary alignment varies according to the type of constant being specified and the presence of a length specification. Some constant types are aligned only to a byte boundary, but the DS instruction can be used to force any type of word boundary alignment for them. This is explained under "DS -- Define Storage." Other constants are aligned at various word boundaries (halfword, fullword, or doubleword) in the absence of a length specification. If length is specified, no boundary alignment occurs for such constants.

Bytes that must be skipped in order to align the field at the proper boundary are not considered to be part of the constant. Thus, the location counter is incremented to reflect the proper boundary (if any incrementing is necessary) before the address value is established. Therefore, the symbol naming the constant will not receive a value that is the location of a skipped byte.

Bytes skipped to align a DC statement are set to zero; bytes skipped to align a DS statement are not set to zero.

Appendix $F$ summarizes, in chart form, the information about constants that is presented in this section.

LITERAL DEFINITIONS: Note that the description of literals in Section 2 referred to the following discussion of the DC operand in reference to the writing of a literal operand. All subsequent operand specifications are applicable to writing literals; the only differences are listed below.

1. The literal is preceded by an equal sign.
2. Multiple constants may not be specified.
3. Unsigned decimal self-defining terms must be used to express the duplication factor and length values.
4. The duplication factor may not be zero.
5. If a reference to the location counter occurs in an address constant that specifies a duplication factor greater than one, the value of the location counter used in each duplication is incremented by the length of the constant; if, however, the reference occurs in a literal address constant, the value remains unchanged throughout duplication.

Examples of literals appear throughout the balance of the DC instruction discussion.

## Operand Subfield 1: Duplication Factor

The duplication factor may be omitted. If specified, it causes the constant(s) to be generated the number of times indicated by the factor. The factor may be specified either by an unsigned decimal self-defining term or by a positive absolute expression that is enclosed by parentheses. All symbols in the expression must be previously defined. A location counter reference may not appear in such an expression. The maximum value permitted for the duplication factor is 65,535.

The duplication factor is applied after the constant is assembled. When more than one constant is specified in a DC operand having a duplication factor, the duplication factor is applied to the constants as a unit, rather than individually. Thus, if a duplication factor of 2 is specified for the constants 1,2 , and 3 , the constants are generated in the order -- $1 \begin{array}{llllll}1 & 2 & 1 & 2 & 3\end{array}$ -- not in the order -- 1122.33.

Note that a duplication factor of zero is permitted and achieves the same result as it would in a DS instruction. A DC instruction with a zero duplication factor will not produce control dictionary entries. See "Forcing Alignment" under "DS -- Define Storage."

Note: If duplication is specified for an address constant containing a location counter reference, the value of the location counter used in each duplication is incremented by the length of the operand. (If the reference occurs in a literal address constant, however, the value remains unchanged.)

## Operand Subfield 2: Type

The type subfield defines the type of constant being specified. From the type specification, the assembler determines how it is to interpret the constant and translate it into the appropriate machine format. The type is specified by a singleletter code as shown in Figure 4.

Further information about these constants is provided in "Operand subfield 4: Constant" below.

## Operand Subfield 3: Length

The length subfield may be omitted. If used, it indicates the length of the specified constant. This is written as Ln, where $\underline{n}$ is either an unsigned decimal self-defining term or a positive absolute expression enclosed by parentheses. Any symbols in the expression must be previously defined. A location counter reference may not appear in such an expression.

The value of $n$ represents the number of bytes of storage that are assembled for the constant. The maximum values permitted for the length of the various types of constants are summarized in Appendix F. This table also indicates the implied length for each type of constant; the implied length is used unless a length subfield is present. A length may be specified for any type of constant. However, no boundary alignment will be provided when a length is given.


Fiqure 4. Type Codes for Constants

## Operand Subfield_4: Constant

This subfield supplies the constant (s) described by the subfields that precede it. A data constant (all types except A) is enclosed by single quotes. An address constant (type A) is enclosed by parentheses. To specify two or more constants in the subfield, the constants must be separated by commas and the entire sequence of constants must be enclosed by the appropriate delimiters (i.e., sinqle quotes or parentheses). Thus, the format for specifying the constant(s) is one of the followinq:

| Sinale | Multiple |
| :--- | :--- |
| Constant | Constants* |
| 'constant | 'constant,....constant' |
| (constant) | (constant,..., constant) |

Fixed-point ( $F$ and $H$ ), floatinq-point (E and D), and address (A) constants are aliqned on the proper boundary, as shown in Appendix $F$, unless a length modifier is specifed. In the presence of a lenqth modifier, no boundary aliqnment is performed. If an operand specifies more than one constant, any necessary alignment applies to the first constant only. Thus, for an operand that provides five fullword constants, the first would be aliqned on a fullword boundary, and the rest would automatically fall on fullword boundaries.

[^2]The total storaqe requirement of an operand is the product of the lenqth multiplied by the number of constants in the operanu, which iñ turn is muitiplied by the duplication factor (if present), plus any bytes skipped for boundary alignment of the first constant.

If an address constant contains a location counter reference, the location counter value that is used is the storaqe address of the first byte the constant will occupy. Thus, if several address constants in the same instruction refer to the location counter, the value of the location counter varies from constant to constant. Similarly, if a sinqle address constant is specified (and it is a location counter reference) with a duplication factor, the constant is duplicated with a varying location counter value.

The following text describes each of the constant types and provides examples.

Character Constant $=-C:$ Any of the valid 256 punch combinations may be desiqnated in a character constant. Only one character constant may be specified per operand. Since multiple constants within an operand are separated by commas, an attempt to specify two character constants would result in interpreting the comma separating them as a character.

Special consideration must be qiven to representing single quotes and ampersands as characters. Each single quote or ampersand desired as a character in the constant must be represented by a pair of sinqle quotes or ampersands. Only one single quote or ampersand appears in storage.

The maximum length of a character con|stant is 256 bytes. No boundary alignment is performed. Each character is translated into one byte. Paired sinqle quotes or paired ampersands count as one character. If no lenqth modifier is given, the size in bytes of the character constant is equal to the number of characters in the constant. If a length is provided, the result varies as follows:

1. If the number of characters in the constant exceeds the specified lenqth, as many riqhtmost bytes as are necessary are dropped.
2. If the number of characters is less than the specified lenqth, the excess rightmost bytes are filled with blanks.

In the following example, the lenqth of FIELD is $12:$


However, in this next example, the length is 15 , and three blanks appear in storaqe to the riqht of the zero:


In the next example, the lenqth of FIELD is 12 , although 13 characters appear in the operand. The two ampersands count as only one byte.


Note that in the next example, a lenath of four has been specified, but there are five characters in the constant.


The generated constant would be:

## ABCDABCDABCD

However, if the length had been specified as six instead of four, the generated constant would have been as shown below (with the spaces between and following the grouped constants being the sixth character) :

## $\operatorname{ABCDE}$ ABCDE ABCDE

Hexadecimal Constant_-_ X: A hexadecimal constant consists of one or more of the hexadecimal digits, which are 0 through 9 and $A$ through $F$. Only one hexadecimal constant may be specified per operand. The maximum length of a hexadecimal constant is 32 bytes (64 hexadecimal digits). No boundary aliqnment is performed.

Constants that contain an even number of hexadecimal digits are translated as one byte per pair of digits. If an odd number of digits is specified, the leftmost byte has the leftmost four bits filled with a hexadecimal zero, while the other four bits contain the odd (first) digit.

If no lenqth modifier is qiven, the implied length of the constant is half the number of hexadecimal diaits in the constant (assuming that a hexadecimal zero is added to an odd number of diqits). If a lenath modifier is qiven, the constant is handled as follows:

1. If the number of hexadecimal diqit pairs exceeds the specified lenqth, as many hexadecimal diqits as necessary are dropped from the left.
2. If the number of hexadecimal diqit pairs is less than the specified length, as many hexadecimal zeros as are necessary are added on the left.

An 8-diqit hexadecimal constant provides a convenient way to set the bit pattern of a full binary word. The constant in the following example would set the first and third bytes of a word to ones (the DS instruction sets the location counter to a fullword boundary):

| \| Name | loperation | 10 per and |
| :---: | :---: | :---: |
| \| | \\| DS | 10 F |
| \|TEST | IDC | $\left.\right\|^{\prime} \mathrm{PF} 00 \mathrm{FF} 00^{\prime}$ |
|  |  |  |

The next example uses a hexadecimal constant as a literal and inserts ones into bits 24 through 31 of reqister 5.


In the following example, the diqit $A$ would be dropped, because five hexadecimal digits are specified for a length of two bytes:


The resulting constant would be $6 F 4 E$, which would occupy the specified 2 bytes. It would then be duplicated 3 times, as requested by the duplication factor. If it had merely been specified as X'A 6 F4E', the resulting constant would have had a hexadecimal zero in the leftmost position, as follows:

OA6F4E

Fixed-Point Constants -- F and H: A fixedpoint constant is written as a signed or unsigned decimal self-defining term. It is assumed that the sign is positive if an unsigned term is specified.

The decimal value is converted to a binary number. If the value of the number exceeds the length specified or implied, the sign is lost, the necessary leftmost bits are truncated to the length of the field, and the value is then assembled into the whole field. Any duplication factor that is present is applied after the constant is assembled. A negative number is carried in two's complement form.

An implied length of 4 bytes is assumed for a fullword (F) and 2 bytes for a halfword (H), and the constant is aligned to the proper fullword or halfword if a length is not specified. However, any length up to and including 8 bytes may be specified for either type of constant by a length modifier, in which case no boundary alignment occurs.

Maximum and minimum values for fixedpoint constants are:

| Length | Maximum | Minimum |
| :---: | :---: | :---: |
| 8 | 263-1 | $-\left(2^{63}-1\right)$ |
| 4 | 231-1 | $-2^{31}$ |
| 2 | 215-1 | $-215$ |
| 1 | 27-1 | -27 |

A field of 3 fullwords is generated from the statement shown below. The value of CONWRD is the address of the leftmost byte of the first word, and the length of the constant is 4, the implied length for a fullword. fixed-point constant. The expression CONWRD +4 could be used to address the second constant (second word) in the fiela.


A constant could be specified as a literal:


Floating-Point Constants -- $E$ and D: A floating-point constant is written as a decimal number, which may be followed by a
decimal exponent, if desired. The number may be an integer, a fraction, or a mixed number (i.e., one with integral and fractional portions). The format of the constant is as follows:

1. The number is written as a signed or unsigned decimal value. The decimal point may be placed before, within, or after the number, or it may be omitted, in which case, it is assumed that the number is an integer. It is also assumed that the sign is positive if an unsigned number is specified.
2. The exponent is optional. If specified, it is written immediately after the number as En, where $\underline{n}$ is a signed or unsigned decimal self-defining term specifying the exponent of the factor 10. The exponent may be in the range from -78 to +75 . If an unsigned exponent is specified, it is assumed that the sign is a plus.

The floating-point constant is converted to a normalized hexadecimal floating-point constant in machine format. Truncation of the fraction is performed according to the specified or implied length, and the number is stored in the proper field. The resulting number will not differ from the exact value by more than one in the rightmost place.

The implied length for a fullword (type E) constant is 4 bytes; the implied length for a double word (type D) constant is 8 bytes. The constant is aligned at the proper word or double word boundary if a length is not specified. However, any length up to and including 8 bytes may be specified for either type of constant by a length modifier, in which case no boundary alignment occurs.

Any of the following statements could be used to specify 46.415 as a positive, full word, floating-point constant; the last is a machine instruction statement with a literal operand.


Each of the following would be generated as double word floating-point constants.


Address Constant -- A: An address constant is specified as an absolute or relocatable expression. (Note that an expression may be single term or multiterm.) The value of the expression is calculated as explained in section 2 with one exception: the maximum value of an absolute expression may be 231-1. The value is then truncated on the left, if necessary, to the specified or implied length of the field and assembled into the rightmost bits of the field. The implied length of an address constant is 4 bytes, and alignment is to a fullword boundary unless a length is specified, in which case no alignment will occur. The length that may be specified depends on the type of expression used for the constant; a length of 1 to 4 bytes may be used for an absolute expression, while a length of only 3 or 4 may be used for a relocatable expression.

Address constants are used for initializing base registers to facilitate the addressing of storage. Furthermore, they provide the means of communicating between control sections of a multisection program. However, storage addressing and control section communication are also dependent on the use of the USING assembler instruction and the loading of the registers. Coding examples that illustrate these considerations are provided in section 3 under the heading "Programming with the USING Instruction."

In the following examples, the field generated from the statement named CONST contains a location counter reference. The value of the location counter will be the address of the first byte allocated to the constant. The second statement shows the same constant specified as a literal (i.e., an address constant literal).


When the location counter reference occurs in a literal, the value of the location counter is the address of the first byte of the instruction in which the literal is used.

## DS -- DEFINE STORAGE

The DS instruction is used to reserve areas of storage and to assign names to those areas. The use of this instruction is the preferred way of symbolically defining storage for work areas, input/output areas, etc. The size of a storage area that can be reserved by using the DS instruction is limited only by the maximum value of the location counter.


The format of the DS operand is identical to that of the DC operand; exactly the same subfields are employed and are written in exactly the same sequence as they are in the DC operand. Although the formats are identical, there are two differences in the specification of subfields, as follows:

1. The specification of data (subfield 4) is optional in a DS operand, but it is mandatory in a DC operand. If the constant is specified, it must be valid.
2. The maximum length that may be specified for character (C) and hexadecimal (X) field types is 65,535 bytes rather than 256 bytes.

If a DS operand specifies a constant in subfield 4, and no length is specified in subfield 3, the assembler determines the length of the data and reserves the appropriate amount of storage. It does not assemble the constant. The ability to specify data and have the assembler calculate the storage area that would be required for such data is a convenience to the programmer. If he knows the general format of the data that will be placed in the storage area during program execution, all the programmer need do is show it as subfield 4 in a DS operand. The assembler then determines the correct amount of storage to be reserved, thus relieving the programmer of length considerations.

If the DS instruction is named by a symbol, its value is the location of the leftmost byte of the reserved area. The length of the field is the length (implied or explicit) of the type of data specified. Any positioning required for aligning the storage area to the proper type of boundary is done before the address value is deter-
mined. Bytes skipped for alignment are not set to zero.

Each field type (e.g., hexadecimal, character, floating-point) is associated with certain characteristics (these are summarized in Appendix F) = The associated characteristics will determine which fieldtype code the programmer selects for the DS operand and what other information he adds, notably a length specification or a duplication factor. For example, the $E$ floating-point field and the F fixed-point field both have an implied length of 4 bytes. The leftmost byte is aligned to a fullword boundary. Thus, either code could be specified if it were desired to reserve 4 bytes of storage aligned to a fullword boundary. To obtain a length of 8 bytes, one could specify either the $E$ or $F$ field type with a length modifier of 8. However, a duplication factor would have to be used to reserve a larger area, because the maximum length specification for either type is 8 bytes. Note also that specifying length would cancel any special boundary alignment.

In contrast, character (C) and hexadecimal (X) fields have an implied length of 1 byte. Either of these codes, if used, would have to be accompanied by a length modifier, unless just 1 byte is to be reserved. Although no alignment occurs. the use of $C$ and $X$ field types permits greater latitude in length specifications, the maximum for either type being 65,535 bytes. (Note that this differs from the maximum for these types in a DC instruction.) Unless a field of 1 byte is desired, either the length must be specified for the $C$ or $X$ field types, or else the data must be specified (as subfield 4), so that the assembler can calculate the length.

To define four 10-byte fields and one 100 -byte field, the respective DS statements might be as follows:


Additional examples of DS statements are shown below:


Note: A DS statement causes the storage area to be reserved but not set to zeros. Assumptions should not be made as to the contents of the reserved area.

## Special Uses of the Duplication Factor

FORCING ALIGNMENT: The location counter can be forced to a double word, fullword, or halfword boundary by using the appropriate field type (e.g., D, F, or H) with a duplication factor of zero. This method may be used to obtain boundary alignment that otherwise would not be provided. For example, the following statements would set the location counter to the next double word boundary and then reserve storage space for a 128-byte field (whose leftmost byte would be on a double word boundary).


DEFINING FIELDS OF AN AREA: A DS instruction with a duplication factor of zero can be used to assign a name to an area of storage without actually reserving the area. Additional DS and/or DC instructions may then be used to reserve the area and assign names to fields within the area (and generate constants if $D C$ is used).

For example, assume that 80-character records are to be read into an area for processing and that each record has the following format:

| Position |  | Meaning |
| ---: | :--- | :--- |
| $11-30$ |  | Payroll Number |
| $31-36$ |  | Employee Name |
| $47-54$ |  | Gross Wages |
| $55-62$ |  | Withholding Tax |

The following example illustrates how DS instructions might be used to assign a name to the record area, then define the fields
of the area and allocate the storage for them.

| \| Name | Operation | \|operand |
| :---: | :---: | :---: |
| \|RDAREA | \|DS | 10CL 80 |
| \| | \|DS | \| CL4 |
| \| PAYNO | \|DS | \| CL6 |
| \| NAME | IDS | \| CL20 |
| \|DATE | \|DS | \| 0CL6 |
| \| DAY | \|DS | \| CL2 |
| \| MONTH | \|DS | \| CL2 |
| \| YEAR | \|DS | \| CL2 |
|  | \| DS | \| CL10 |
| \| GROSS | \|DS | \| CL8 |
| \|FEDTAX | \|DS | \|CL8 |
|  | \|DS | \|CL18 |

Note that the first statement names the entire area by defining the symbol RDAREA; but does not reserve any storage. Similarly, the fifth statement names a 6-byte area by defining the symbol DATE; the three subsequent statements actually define the fields of DATE and allocate storage for them. The second, ninth, and last statements are used for spacing purposes and, therefore, are not named.

## CCW -- DEFINE CHANNEL COMMAND WORD

The CCW instruction provides a convenient way to define and generate an 8 byte channel command word aligned at a doubleword boundary. The internal machine format of a channel command word is shown in Table 2. The format of the CCW instruction statement is:

| \| Name | \|operation | \| Operand |
| :---: | :---: | :---: |
| \|A symbol| | CCW | \| Four operands, |
| \|or blank| |  | \|separated by commas, |
|  |  | \|specifying the con- |
| \| |  | \|tents of the channel |
| I |  | \| command word in the |
| 1 |  | \|format described in |
| ! |  | \|the following text |

All four operands must appear. They are written, from left to right, as follows:

1. An absolute term that specifies the command code. The value of this term is right-justified in byte 1.
2. An absolute or relocatable expression specifying the data address. The value of this expression is rightjustified in bytes 2 through 4.
3. An absolute term that specifies the flags for bits 32 through 36 and zeros for bits 37 through 39. The value of this term is right-justified in byte 5. (Byte 6 is set to zero.)
4. An absolute term that specifies the count. The value of this term is right justified in bytes 7 and 8.

The following is an example of a CCW statement:


Note that the form of the third operand sets bits 37 through 39 to zero, as required. The bit pattern of this operand is as follows:

$$
\frac{32-35}{0100} \quad \frac{36-39}{1000}
$$

If there is a symbol in the name field of the CCW instruction, it is assigned the address value of the leftmost byte of the channel command word.

Table 2. Channel Command Word

| Byte | \|Bits | \|Usage |
| :--- | :--- | :--- |
| 1 | $0-7$ | \| Command code |
| $2-4$ | $8-31$ | \| Data address |
| 5 | $32-36$ | \|Flags |
| 6 | $37-39$ | \| Must be zero |
| 6 | $40-47$ | \|Set to zero |
| $7-8$ | $48-63$ | \|Count |

## LISTING CONTROL INSTRUCTIONS

The listing control instructions are used to identify the program listing and assembly output cards, to provide blank lines in the program listing, and to designate how much detail is to be included in the program listing. In no case are instructions or constants generated in the object program. Listing control statements with the exception of PRINT are not printed in the listing.

## TITLE -- IDENTIFY ASSEMBLY OUTPUT

The TITLE instruction enables the programmer to identify the program listing and
assembly output cards. The format of the TITLE instruction statement is as follows:


The name field may contain an $I D$ field of from one to four alphabetic or numeric characters in any combination. The contents of the ID field are punched into columns 73 through 76 of all the output cards for the program except those produced by the REPRO assembler instruction. Only the first TITLE statement in a program may make use of the ID field. The ID field of all subsequent TITLE statements must be blank.

The operand field may contain up to 62 characters enclosed in single quotes. Special consideration must be given to representing single quotes and ampersands as characters. Each single quote or ampersand desired as a character in the constant must be represented by a pair of single quotes or ampersands. Only one single quote or ampersand appears in storage. The contents of the operand field are printed at the top of each page of the program listing.

A program may contain more than one TITLE statement. Each TITLE statement provides the heading for subsequent pages in the program listing, until another TITLE statement is encountered. Each TITLE statement causes the listing to be advanced to a new page (before the heading is printed).

For. example, if the following statement is the first TITLE statement to appear in a program:

then PGM1 is punched into all of the output cards (columns 73 through 76) and FIRST HEADING appears at the top of each subsequent page.

If the following statement occurs later in the same program:

then, PGM1 is still punched into the output cards, but í NEW HEADING appears at the top of each subsequent page.

Note: A title card with a non-blank name field must be used if the output deck will at any time be processed by the update feature of the assembler. In conjunction with the sequence numbers punched automatically in columns 77 through 80, a 4-character ID provides 8-digit serialization in columns 73 through 80, as required for an update operation.

## EJECT -- START NEW PAGE

The EJECT instruction causes the next line of the listing to appear at the top of a new page. This instruction provides a convenient way to separate routines in the program listing. The format of the EJECT instruction statement is as follows:

| Name | Operation | \|operand |
| :---: | :---: | :---: |
| Blank | \| EJECT | \|Must be blank |

If the line before the EJECT statement would have been the last line on a page anyway, the EJECT statement has no effect. Two EJECT statements may be used in succession to obtain a blank page. A TITLE instruction followed immediately by an EJECT instruction will produce a page with nothing but the operand entry of the TITLE instruction. Text following the EJECT instruction will begin at the top of the next page.

SPACE -- SPACE LISTING

The SPACE instruction is used to insert one or more blank lines in the listing. The format of the SPACE instruction statement is as follows:


A decimal value is used to specify the number of blank lines to be inserted in the program listing. A blank operand causes one blank line to be inserted. If the operand value exceeds the number of lines remaining on the listing page, the statement will have the same effect as an EJECT statement.

PRINT -- PRINT OPTIONAL DATA

The PRINT instruction is used to control printing of the program listing. The format of the PRINT instruction statement is:


One or both of the following operands are used:

## 1. ON - A listing is printed. <br> OFF - No listing is printed.

2. DATA - Constants are printed out in full in the listing.

NODATA - Only the leftmost 8 bytes are printed in the listing.

A program may contain any number of PRINT statements. The PRINT statement controls the printing of the program listing until another PRINT statement is encountered.

Until the first PRINT statement (if any) is encountered, the following is assumed:


For example, if the statement:

appears in a program, 256 bytes of zeros are assembled. If the statement:

is the last PRINT statement to appear before the DC statement, all 256 bytes of zeros are printed in the program listing. However, if the following statement is the last PRINT statement to appear before the DC statement, only 8 bytes of zeros are printed in the program listing.


Whenever an operand is omitted, it is assumed to be unchanged and continues according to its last specification.

If the OFF operand is used, no data will be printed even though a DATA operand is specified. Thus, with the following statement nothing would be printed.


## PROGRAM CONTROL INSTRUCTIONS

The program control instructions are used to specify the end of an assembly, to set the location counter to a value or word boundary, to specify the placement of literals in storage, to indicate statement format, and to punch a card. Except for the CNOP instruction, none of these assembler instructions generate instructions or constants in the object program.

ICTL -- INPUT FORMAT CONTROL

The ICTL instruction allows the programmer to alter the normal format of his source program statements. The ICTL statement must precede all other statements in the source program and may be used only once. The format of the ICTL instruction statement is as follows:


The operand specifies the begin column of the source statement.

If no ICTL statement is used in the source program, the assembler assumes that column 1 is the begin column.

## REPRO -- REPRODUCE FOLLOWING CARD

The REPRO assembler-instruction causes data on the following statement line to be punched into a card; the data is not processed. Neither a sequence number nor the identification is punched on the card. One REPRO instruction produces one punched card.

A REPRO statement that occurs before all program sectioning and linking instructions and before any assembler language instruction that may either affect or depend upon the setting of the location counter causes the assembler to punch a card that precedes all other cards of the object deck. A REPRO statement that occurs after any of the program sectioning and linking instructions has been encountered causes the assembler to punch a card that follows the object cards produced for all of the program sectioning and linking instructions. (The program sectioning and linking instructions are: START, CSECT, DSECT, COM, ENTRY, and EXTRN.)

```
The format of the REPRO instruction statement is:
```



The line to be reproduced may contain any combination of valid characters, starting in column 1 and continuing through column 72 of the line. Column 1 of the line corresponds to column 1 of the card to be punched.

ORG -- SET LOCATION COUNTER

The ORG instruction is used to alter the setting of the location counter for the current control section. The format of the ORG instruction statement is:


Any symbols in the expression must have been previously defined. An unpaired relocatable symbol must be defined in the same control section in which the ORG statement appears.

The location counter is set to the value of the expression in the operand. If the operand is omitted, the location counter is set to the next available (unused) location for that control section.

An ORG statement must not be used to specify a location below the beginning of the control section in which it appears. For example, the following statement is invalid if it appears less than 500 bytes from the beginning of the current control section.


If it is desired to reset the location counter to the highest location yet assigned (in the control section), the following statement would be used:


If previous ORG statements have reduced the location counter for the purpose of
redefining a portion of the current control section, an ORG statement with an omitted operand can then be used to terminate the effects of such statements and restore the location counter to its highest setting.

LTORG -- BEGIN LITERAL POOL

The LTORG instruction causes all literals since the previous LTORG (or start of the current control section) to be assembled at appropriate boundaries starting at the first double-word boundary following the LTORG statement. If no such literals exist, alignment of the next instruction (which is not a LTORG instruction) will occur. Bytes skipped are not set to zero. The format of the LTORG instruction statement is:


The symbol represents the address of the first byte of the literal pool. The LTORG statement forces all literals in a control section to be generated as a part of that control section. A LTORG statement must appear after the last reference to any literal in a control section.

## Duplicate Literals

If duplicate literals occur within the range controlled by one LTORG statement, only one literal is stored. Literals are considered duplicates only if their specifications are identical. A literal will be stored, even if it appears to duplicate another literal, if it is an A-type address constant containing any reference to the location counter.

The following examples illustrate how the assembler stores pairs of literals, if the placement of each pair is controlled by the same LTORG statement.
$X^{\prime} F 0^{\prime}$
$C^{\prime} 0^{\prime}$
XL3 ${ }^{\prime} 0^{\prime \prime}$
Both are stored
HL3' ${ }^{\prime \prime}$
A(*+4) Both are stored
A(*+4)
$X^{\prime} \mathrm{FFFF}^{\prime}$
Identical; the first is stored
$X^{\prime}$ FFFF'
$X^{\prime} \mathrm{FF}^{\prime}$
Both are stored
XL1 ${ }^{\prime} F F^{\prime}$

CNOP -- CONDITIONAL NO OPERATION

The CNOP instruction allows the programmer to align an instruction at a specific halfword boundary. If any bytes must be skipped in order to align the instruction properly, the assembler ensures an unbroken instruction flow by generating no-operation instructions. This facility is useful in creating calling sequences that consist of a linkage to a subroutine followed by parameters.

The CNOP instruction ensures the alignment of the location counter setting to a halfword, fullword, or double word boundary. If the location counter is already properly aligned, the CNOP instruction has no effect. If the specified alignment requires the location counter to be incremented, one to three no-operation instructions are generated, each of which uses 2 bytes.


Figure 5. CNOP Alignment

The format of the CNOP instruction statement is as follows:


Any symbols used in the expressions in the operand field must previously have been defined. A location counter reference may not appear in such an expression.

Operand b specifies at which byte in a fullword or double word the location counter is to be set; $\underline{b}$ can be $0,2,4$, or 6 . operand $w$ specifies whether byte $b$ is in a fullword ( $\mathbf{w}=4$ ) or double word ( $w=8$ ). The following pairs of $\underline{b}$ and $\underline{w}$ are valid:

| $\mathrm{b}, \mathrm{w}$ | Specifies |
| :--- | :--- |
| 0,4 | Beginning of a fullword |
| 2,4 | Middle of a fullword |
| 0,8 | Beginning of a double word |
| 2,8 | Second halfword of a double word |
| 4,8 | Middle (third halfword) of a |
| double word |  |
| 6,8 | Fourth halfword of a double word |

Figure 5 shows the position in a double word that each of these pairs specifies. Note that both 0,4 and 2,4 specify two locations in a double word.

Assume that the location counter is currently aligned at a double word boundary. Then the CNOP instruction in the following sequence has no effect; it is merely printed in the assembly listing:


However, the following sequence:



After the BALR instruction is generated, the location counter is at a double word boundary, thereby ensuring an unbroken instruction flow.

END -- END ASSEMBLY

The END instruction terminates the assembly of a program. It may also designate a point in the program or in a separately assembled program to which control may be transferred after the program is loaded. The END instruction must always be the last statement in the source program. A literal may not be used.

The typical form of the END instruction statement is as follows:


The operand specifies the point to which control may be transferred when loading is complete. This point is usually the first machine instruction in the program, as shown in the following sequence. If the operand field is blank, control is automatically transferred to the first byte of the first control section in the assembly.

| \| Name | Operation | Operand |
| :---: | :---: | :---: |
| \| NAME | \|CSECT |  |
| \| AREA | \| DS | 150F |
| \| BEGIN | \| BALR | 12,0 |
| I | \| USING | 1*, 2 |
| \| | 1 | \| |
| I | I | ! |
| 1 | 1 |  |
| \| | \| END | \| BEGIIN |

The conditional assembly instructions allow the programmer to bypass source statements during an assembly, depending on the values assigned to variable symbols.

| There are 4 conditional assembly |
| :---: |
| instructions: |
| SETA AIF AGO ANOP |

Note: Other System/360 programming support system assembler languages employ an additional conditional assembly instruction (the LCLA instruction) for the definition of variable symbols. Since the Model 44 Programming System Assembler Language does not require explicit definition of variable symbols beyond their appearance in the name field of a SETA instruction, LCLA statements are not required in the language. To enable an additional degree of compatibility, however, LCLA statements are treated as comments by this assembler.

The SETA instruction is used to define a variable symbol and assign an arithmetic value to it.

The AIF, AGO, and ANOP instructions are used (in conjunction with sequence symbols) to indicate which statements are to be processed by the assembler. The programmer can test the values assigned to variable symbols, thereby determining which statements are to be processed.

An example illustrating the use of conditional assembly instructions is included at the end of this section.

## VARIABLE SYMBOLS

A variable symbol is a type of symbol that is assigned different values by the programmer. A variable symbol is written as an ampersand followed by from one to seven letters and/or digits, the first of which must be a letter. A variable symbol may be used in any operand where a selfdefining term is allowed.

## SETA -- SET ARITHMETIC

The SETA instruction is used to assign an arithmetic value to a variable symbol. A variable symbol is defined when it
appears in the name field of a SETA instruction. The format of this instruction is:


The expression in the operand field is evaluated as a signed 24-bit arithmetic value which is assigned to the variable symbol in the name field. The minimum and maximum allowable values of the expression 1 are $-2^{23}$ and $+2^{23-1}$, respectively.

The expression may consist of one term or an arithmetic combination of up to three terms. The terms may be either selfdefining terms or variable symbols. (Self-defining terms are described in Section 1 of this publication.)

The arithmetic operators that may be used to combine the terms of an expression are + (addition), - (subtraction), * (multiplication), and / (division).

An expression may not contain two terms or two operators in succession, nor may it begin with an operator.

The following are valid operand fields of SETA instructions:

```
GAREA+X'2D' EHERE-EEXIT
EBETA*10 29
```

The following are invalid operand fields of SETA instructions:

$$
\begin{array}{ll}
\text { EAREAX'C' } & \text { (two terms in succession) } \\
\text { \&FIELD+-3 } & \begin{array}{l}
\text { (two operators in } \\
\text { succession) }
\end{array} \\
- \text { EDELTA*2 } & \begin{array}{l}
\text { (begins with an operator) } \\
\text { (begins with an operator: } \\
\text { two operators in }
\end{array} \\
\text { succession) }
\end{array}
$$

## Evaluation of Arithnetic Expressions

The procedure used to evaluate the arithmetic expression in the operand field of a SETA instruction is the same as that used to evaluate arithmetic expressions in assembler language statements. The only
difference between the two types of arithmetic expressions is the terms that are allowea iñ each expression.

The following evaluation procedure is used:

1. Each term is given its numerical value.
2. If a variable symbol is used, the arithmetic value assigned to it is substituted. If no arithmetic value has been assigned to the variable symbol, it is assumed the value is zero.
3. The arithmetic operations are performed moving from left to right. However, multiplication and/or division are performed before addition and subtraction.
4. The computed result is the value assigned to the variable symbol in the name field.

One level of parentheses may be used in a SETA operand. Each term enclosed by parentheses counts toward the maximum of three terms.

The following are examples of SETA instruction operand fields that contain parenthesized terms:

$$
\begin{aligned}
& (\text { \&HERE }+32) * 29 \\
& \text { GEXIT/ }(\text { EENTRY-4) } \\
& \text { EBETA }+(\text { EENTRY* })
\end{aligned}
$$

The parenthesized portion of an arithmetic expression is evaluated before the other term in the expression is evaluated.

## LOGICAL EXPRESSIONS

Logical expressions enable the programmer to test the values assigned to variable symbols. A logical expression is used in the operand field of an AIF instruction and is evaluated to determine if it is true or false.

A logical expression consists of an arithmetic relation enclosed by parentheses. No blanks should appear between the enclosing parentheses and the first or last character of the arithmetic relation.

An arithmetic relation consists of two arithmetic expressions connected by a relational operator. The relational operator must be immediately preceded and followed by at least one blank. The relational operators are $E Q$ (equal), $N E$ (not equal),

LT (less than), GT (greater than), LE (less than or equal), and GE (greater than or equai).

The arithmetic expressions that may be used in an arithmetic relation are limited to those expressions that are valid in the operand field of a SETA instruction.

The following are valid logical expressions:

```
(7*(&ALPHA+6) EQ &GAMMA)
(EAREA+2 GT 29)
(\varepsilonXYZ NE &P12*(&A+6))
```

The following are invalid logical expressions:

EBEQ EA (not enclosed in ( $\varepsilon P$ EQ \&B 6) (two terms in succession) ( \&A EQ 5 ) (blank following left parenthesis and preceding right parenthesis)

## SEQUENCE SYMBOLS

Sequence symbols provide the programmer with the ability to vary the sequence in which statements are processed by the assembler.

A sequence symool may be used in the name field of any statement (except an ICTL statement) that does not require an ordināy symuol.

A sequence symbol is used in the operand field of an AIF or AGO statement to refer to the statement named by the sequence symbol.

A sequence symbol consists of a period followed by one to seven letters and/or digits, the first of which must be a letter.

The following are valid sequence symbols:

| . READER | .A23456 |
| :--- | :--- |
| .LOOP 2 | .X4F2 |
| . N | .S4 |

The following are invalid sequence symbols:
\(\left.$$
\begin{array}{ll}\text { CARDAREA }\end{array}
$$ \begin{array}{l}(first character is not <br>

a period)\end{array}\right]\)| (first character after |
| :---: |
| period is not a letter) |


| . BCD\%84 | (contains a special <br> character other than <br> initial period) |
| :---: | :---: |
| .IN AREA $\left.\begin{array}{c}\text { (contains a special } \\ \text { character, i.e. blank, } \\ \text { other than initial period) }\end{array}\right)$ |  |

## AIF -- CONDITIONAL BRANCH

The AIF instruction is used to conditionally alter the sequence in which source program statements are processed by the assembler. The assembler assigns a maximum count of 4096 AIF and AGO branches that may be executed in the source program. The format of this instruction is:


Any logical expression may be used in the operand field of an AIF instruction. The sequence symbol in the operand field must immediately follow the closing parenthesis of the logical expression.

The logical expression in the operand field is evaluated to determine if it is true or false. If the expression is false, the next sequential statement is processed by the assembler. If the expression is true, the statement named by the sequence symbol in the operand field is the next statement processed by the assembler. This statement must not precede the AIF instruction.

The following are valid operand fields of AIF instructions:

```
(EAREA+X'2D' GT 29).READER
((32-&HERE)*4 GT 48).THERE
```

The following are invalid operand fields of AIF instructions:

| (EHERE NE 6) |  | (no sequence symbol) |
| :---: | :---: | :---: |
| . X 4 F 2 |  | (no logical expression) |
| ( $6 \mathrm{~N}+4 \mathrm{GT}$ EL) | . PASS | (blanks between logical |
|  |  | expression and sequence symbol) |

AGO -- UNCONDITIONAL BRANCH

The AGO instruction is used to unconditionally alter the sequence in which source
program statements are processed by the assembler. The assembler assigns a maximum count of 4096 AIF and AGO branches that may be executed in the source program. The format of this instruction is:


The statement named by the sequence symbol in the operand field is the next statement processed by the assembler. This statement must not precede the AGO instruction.

## ANOP -- ASSEMBLY NO OPERATION

The ANOP instruction facilitates conditional and unconditional branching to statements named by symbols or variable symbols. The format of this instruction is:


If the programmer wants to use an AIF or AGO instruction to branch to another statement, he must place a sequence symbol in the name field of the statement to which he wants to branch. However, if the programmer has already entered a symbol or variable symbol in the name field of that statement, he cannot place a sequence symbol in the name field. Instead, the programmer must place an ANOP instruction before the statement and then branch to the ANOP instruction. This has the same effect as branching to the statement immediately after the ANOP instruction.

## USING CONDITIONAL ASSEMBLY INSTRUCTIONS

The following coding is an example of the use of conditional assembly instructions within a program.


The SETA instructions define the variable symbols EALPHA and EBETA and assign to them the arithmetic values 3 and 4; respectively. The AIF instruction tests these values to determine the next statement to be processed by the assembler. The logical
expression ( $\mathcal{A L P H A} E Q$ GBETA) is evaluated and found to be false: EALPHA is not equal in value to EBETA. Therefore, the assembler continues with the next sequential statement. Statements 1 and 2 are processed. The AGO instruction then causes the assembler unconditionally to bypass statements until it encounters a statement with the sequence symbol .OTHER in the name field. Thus, statements 4, 5, and 6 are bypassed; statement 7 is the next statement processed.

For a subsequent assembly, the SETA instructions can be replaced so that the values of EALPHA and EBETA are equal. Under these circumstances the logical expression in the AIF instruction is true and the next statement to be processed is the statement with the sequence symbol . SKIP in the name field. Thus, the assembler bypasses statements 1, 2, and 3 and processes statement 4. This is an ANOP instruction, required only because the statement following it has an ordinary symbol in the name field. Statements 5, 6, and 7 are then processed by the assembler.

The assembler has an update feature which allows the user to update a serialized, card-image data set by inserting, replacing, or deleting one or more records. The update feature can also be used to serialize a data set for use in subsequent update operations, or to serialize while updating. If the data set consists of source language statements written in assembler language, the assembler may be instructed to assemble the updated source data set concurrently with the update process.

## INPUT/OUTPUT CONSIDERATIONS

The update feature may be used to modify any EBCDIC data set that consists of cards or card images, either blocked or unblocked, provided columns 73 through 80 of each card image are available for or contain a valid serial number. A valid serial number consists of 8 alphameric, non-blank characters. It usually contains a low-order, numeric field sufficient to accommodate numerical sequencing (in increments of 10 ) of the entire data set.

An object deck produced by the assembler can be updated only if produced during an assembly that used a TITLE instruction with a non-blank name field. The 4-character name field of the TITLE instruction is reproduced in columns 73 through 76 of the object program output and, in conjunction with sequence numbers punched automatically in columns 77 through 80, provides a valid serial number for a subsequent update operation.

The update input generally consists of two data sets: an old data set, and an edit data set. Updating of the old data set proceeds under the control of the edit data set.

The old data set consists of card images in sequence by serial number. The edit data set consists of additional card-image data and control statements, and must also be in sequence by serial number, except as noted below. Certain update modes of operation do not require an old data set.

The update output consists of an updated data set and an update listing. (Production of the updated data set is optional for an update-and-assemble run.) The update listing is separate from the
program listing produced by the assembler. It describes the results of the update operation, and includes diagnostic error messages for any errors encountered while updating.

The user specifies a mode of operation appropriate to his input/output requirements in a control card supplied at assembler execution time. The various options and associated input/output assignments are discussed in the publication IBM System/360 Model 44 Programming System: Guide to System Use, Form c28-6812.

## UPDATE OPERATION

## PROCEDURE

The update feature uses the EBCDIC collating sequence to compare the serial numbers of the current records of the old data set and the edit data set. As long as the edit serial is greater, card images are passed from the old data set to the output data set. As soon as the edit serial is equal to or less than the old serial, the edit record is acted upon, as follows:

1. If the edit record is an update instruction, the instruction is performed. The specific actions relating to each update instruction are described under "Update Instructions," below. Generally, their performance involves a repositioning of the old data set, subsequent to which the edit data set is advanced to the next record, and a new comparison is initiated as above.
2. If the edit record is not an update instruction, the edit card image is inserted into the output data set and the edit data set is advanced to the next record. If the comparison yielded an equality, the old data set is also advanced to the next record, thus effecting replacement of the old record with the corresponding edit record. A new comparison is then initiated, and processing continues as above.

Unserialized card images (columns 73 through 80 are all blank) nay appear in either data set at any time. They are considered to have the lowest value, and
are processed immediately upon being encountered. If, however, blank serialization occurs simultaneously in both data sets, the edit record is acted upon.

Note: The following considerations apply during an update-and-assemble operation in which both an old data set and an edit data set are employed. Because the assembier immediately processes an end-of-data (/*) statement with blank serialization encountered in the edit data set, it will not process records remaining in the old data set. To enable processing of the old data set in its entirety, it is necessary to either: (1) serialize the end-of-data statement in such a way that its serial number is greater than the serial number of the last record in the old data set, or (2) ensure that the record immediately preceding the end-of-data statement in the edit data set acts, as required, upon or through the last record in the old set. During an update-only operation, in which job definition statements are not recognized as such, the only requirement is that the end-ofdata statement follow the special control statement (the ENDUP instruction) required to terminate an update-only operation. (For a complete discussion of the end-ofdata statement, see the publication, IBM System/360 Model 44 Programming System: Guide to System Use, Form C28-6812.)

## INSERTION AND REPLACEMENT

Card images may be inserted or replaced without using update instructions. New card images are placed in the edit data set and processed as follows:

- A serialized card image in the edit data set replaces a card image with matching serialization in the old data set.
- A card image with unmatching serialization is inserted in sequence.
- Unserialized card images are inserted immediately upon being encountered.

If only the first of a group of cards in the edit data set contains a serial number, the entire group is inserted.

Card images inserted or replaced are noted as such in the update listing.

## UPDATE INSTRUCTIONS

Functions other than insertion and replacement are specified by update instructions inserted by the user in the edit data set. The OMIT and CPYTO instructions allow a specified segment of the old data set to be omitted from or copied to the new data set. The SKPTO and REWND instructions enable a specified repositioning of the old data set, either forward to a specified serial number, or backward to the first record.

Two additional instructions may be used. The NUM instruction causes reserialization of the new data set. The ENDUP instruction terminates an update-only operation.

A detailed description of each update instruction follows. The following considerations apply:

1. Update instructions in the edit data set function as control statements to the update feature. They are not inserted into the new data set, which may, subsequently or concurrently, be processed by the language translation facility of the assembler. (If, however, an update instruction appears in the old data set, it will not be recognized as such, will not affect the update process, and will be inserted into the new data set unless replaced or deleted by the edit data set.) Update instructions are invalid as input during an assemble-only run.
2. The OMIT, SKPTO, CPYTO, and REWIND instructions refer to an old data set. If encountered during a run in which an old data set is not used, they are flagged as invalid in the update listing and ignored.
3. The NUM and ENDUP instructions may use the serial field to refer to an old data set. Hence, the serial field should be blank when these instructions are used during a run in which an old data set is not used. If, during such a run, a NUM or an ENDUP instruction with a nonblank serial field is encountered, the serial number is ignored, and the instruction is acted upon at once.

## NUM Instruction

The NUM instruction is used to reserialize columns 73 through 80 of the records in the new data set. The format of the instruction is:


The operand field contains 8 characters that specify the starting value of the new serial number. The number in the serial field identifies the record in the old data set at which reserialization is to begin.

When a NUM instruction is encountered, records are read from the old data set and written into the new data set until a serial number is encountered that matches the serial number of the instruction. The operation is then initiated, and continues until a new NUM instruction is encountered or until the last record has been numbered. A NUM instruction with a blank operand field may be used to terminate reserialization of the new data set.

If the serial field of a NUM instruction is blank, reserialization is initiated or terminated immediately.

Reserialization from record to record is done in increments of ten. The rightmost portion of the serial number specified in the operand field must contain enough numeric positions to prevent incrementation from affecting a non-numeric character.

As an example, consider the following statement:


Starting with the record in the old data that has serial number ABC24710, records in the new data set will be serialized IJK00000, IJK00010, IJK00020, etc.

## OMIT Instruction

The OMIT instruction causes deletion of one or more records appearing in the old data set. Deleted records are noted as such in the update listing. The format of the OMIT instruction is:


The serial field contains the serial number at which deletion begins. The operand field contains the serial number at which deletion ends. Those records that have serial numbers that are equal to or between these two serial numbers will be deleted.

When an OMIT instruction is encountered, records are read from the old data set into the new data set until a serial number is encountered that is greater than or equal to the serial number of the instruction. The operation is then initiated, and continues until a serial number is encountered that is greater than or equal to the operand field of the instruction.

If the operand field is blank, the operation is terminated subsequent to the deletion of the single record specified.

If the serial field is blank, the statement is flagged as erroneous in the update listing and ignored.

## SKPTO Instruction

The SKPTO instruction causes the bypassing of one or more records that appear in the old data set. Bypassed records do not appear in the new data set, are not assembled, and do not appear in the update listing.

The format of the SKPTO instruction is:


The serial field contains the serial number of the record immediately preceding the first record to be bypassed. The operand field contains the serial number of the record immediately following the last record to be bypassed. Note that the records whose serial numbers are specified are not themselves bypassed.

When a SKPTO instruction is encountered, records are read from the old data set and written into the new data set until a
serial number is encountered that matches the serial number of the instruction. The operation is then initiated, and continues until a serial number is encountered that matches the operand field of the instruction.

If the serial field is blank, the operation is initiated immediately.

## CPYTO Instruction

The CPYTO instruction causes one or more records to be copied from the old data set into the new data set. Copied records are not assembled and do not appear in the update listing. The format of the CPYTO instruction is:


The serial field contains the serial number of the record immediately preceding the first record to be copied. The operand field contains the serial number of the record immediately following the last record to be copied. Note that the records whose serial numbers are specified are not themselves copied.

When a CPYTO instruction is encountered, records are read from the old data set and written into the new data set until a serial number is encountered that matches the serial number of the instruction. The operation is then initiated, and continues until a serial number is encountered that matches the operand field of the instruction.

If the serial field is blank, the operation is initiated immediately.

## REWND Instruction

The REWND instruction causes the unit holding the old data set to be repositioned to the first record in the data set. In conjunction with the SKPTO instruction, it allows the user to rearrange major segments of the old data set.

The format of the REWND instruction is:


When a REWND instruction is encountered, records are read from the old data set and written into the new data set until a serial number is encountered that matches the serial number of the instruction. The corresponding record is then inserted into the new data set, subsequent to which the repositioning of the old data set is performed.

If the serial field of the instruction is blank, repositioning takes place immediately.

## ENDUP Instruction

The ENDUP instruction is required to terminate an update-only operation. ENDUP should not be issued when the data set is being both updated and assembled. The format of the ENDUP instruction is:


When an ENDUP instruction is encountered, records are read from the old data set and written into the new data set until a serial number is encountered that matches the serial number of the instruction. The corresponding record is then inserted into the new data set, subsequent to which the update operation is terminated.

If the serial field of the instruction is blank, updating is terminated immediately.

Note: An ENDUP instruction must appear in the edit data set during an update-only run. It is the only input that delimits the edit data set. If the assembler does not encounter an ENDUP instruction, it will continue to process the edit data set until it encounters end-of-file.

## SEQUENCE CHECK ING

During an update operation that employs both an old data set and an edit data set, the serialization of the edit data set is checked for proper sequencing. In general, a nonblank serial number must be higher than the preceding nonblank serial number in the edit data set. Blank serialization may, however, occur at any time, and does not constitute a sequence error.

Two exceptions apply. The serial number of a NUM instruction may be equal to (but not greater than) the serial number of the next record in the edit data set. The serial number of the edit record following a REWND instruction may assume any value.

An additional requirement applies to those update instructions (viz.. OMIT, SKPTO, and CPYTO) that specify a range of operation by using the serial field and the operand field, respectively, to identify the initial and terminal extents of the range. For such instructions, a nonblank operand field participates in the sequence check, and must conform to one of the following two requirements:

- A nonblank operand of an OMIT instruction must be greater than the serial number of the instruction and smaller than the serial number of the next record in the edit data set.
- The operand field of a SKPTO or CPYTO instruction must be greater than the serial number of the instruction and smaller than or equal to the serial number of the next record in the edit data set.

If a serialization error is encountered, an error message is written in the update listing. Further action depends on the type of statement being processed, as follows:

- If the record is not an update instruction, it is inserted immediately into the new data set.
- If the record is an OMIT instruction, the instruction is ignored.
- If the record is any other update instruction, records are read from the old data set and written into the new data set until a matching serial number is found.

Note that update does not recognize sequence errors in the old data set. Their occurrence may cause unpredictable results.

## EXAMPLES OF UPDATE OPERATION

Each of the following examples desribes an update operation by specifying two input data sets and the resulting output data set.

Example 1: The old data set is as follows:

| \|columns 1 | 1-72 | \| Columns |
| :---: | :---: | :---: |
| Old Data | 1 | \| Ex100010 |
| jold Data | 2 | \| Ex100020 |
| jold Data | 3 | \|EX100030 |
| Old Data | 4 | \|Ex100040 |
| Old Data | 5 | \|EX100050 |
| Old Data | 6 | \| Ex100060 |
| fold Data | 7 | \| Ex100070 |
| fold Data | 8 | \|EX100080 |
| \|old Data | 9 | \| EX100090 |
| fold Data | 10 | \| Ex100100 |
| \|old Data | 11 | \| EX100110 |

The user wishes to insert Edit Data 1 between Old Data 2 and Old Data 3; to delete Old Data 5 and replace it with Edit Data 2, 3, and 4; to omit Old Data 7. 8, and 9; and to insert Edit Data 5 between Old Data 10 and Old Data 11. The following edit data set will perform these modifications, during an update-only run:

| \| Columns 1-72 | \| Columns 73-80| |
| :---: | :---: |
| \|Edit Data 1 | \|Ex100025 |
| \|Edit Data 2 | EX100050 |
| \|Edit Data 3 | (Blank) |
| \|Edit Data 4 | \|(Blank) |
| 1 OMIT EX100090 | \| EX100070 |
| \|Edit Data 5 | \| Ex100105 |
| ENDUP | EX100110 | operation is as follows:


| \| Columns 1-72 | \| Columns 73-80| |
| :---: | :---: |
| \|old Data 1 | \| Ex100010 |
| \|Old Data 2 | \| EX100020 |
| Edit Data 1 | \| Ex100025 |
| \|old Data 3 | \| EX100030 |
| \|Old Data 4 | \| Ex100040 |
| Edit Data 2 | \| Ex100050 |
| Edit Data 3 | \| (Blank) |
| Edit Data 4 | ( Blank) |
| \|Old Data 6 | \| Ex100060 |
| \|old Data 10 | \| Ex100100 |
| \|Edit Data 5 | \| Ex100105 |
| \|Old Data 11 | \| EX100110 |

Example 2: The user wishes to rearrange and reserialize the following old data set:

| \|columns 1-72 | \|Columns 73-80| |
| :---: | :---: |
| Old Data 1 | \| EX2A0010 |
| Old Data 2 | \| Ex2A0020 |
| Old Data 3 | -Ex2A0030 |
| 101d Data 8 | \| EX2A0040 |
| fold Data 9 | \| EX2A0050 |
| Old Data 10 | \|EX2A0060 |
| lold Data 7 | - EX2A0070 |
| jold Data 4 | \|EX2A0080 |
| Old Data 5 | - Ex2A0090 |
| 1Old Data 6 | \| EX2A0100 |
| \|Old Data 11 | \| EX2A0110 |

In order to accomplish this, the following edit data set is employed, during an update-only run.

| \| Columns | 1-72 | \| Columns 7 |
| :---: | :---: | :---: |
| \| NUM | EX2B0010 | ( Blank) |
| \| SKPTO | EX2A0080 | \| EX2A0030 |
| 1 REWND |  | \|EX2A0100 |
| \| SKPTO | EX2A0070 | ( ${ }^{\text {Blank) }}$ |
| \| REWND |  | \| EX2A0070 |
| \| SKPTO | EX2A0040 | \|(Blank) |
| \| SKPTO | EX2A0110 | \| EX2A0060 |
| \| ENDUP |  | \|EX2A0110 |

The resulting new data set is as follows:

| Columns 1-72 | \| Columns 73-80| |
| :---: | :---: |
| \|Old Data 1 | \|Ex2B0010 |
| \|Old Data 2 | \| EX2B0020 |
| fold Data 3 | \| EX2B0030 |
| \|Old Data 4 | \|EX2B0040 |
| \|Old Data 5 | \| EX2B0050 |
| IOld Data 6 | \|EX2B0060 |
| \|Old Data 7 | \| Ex2B0070 |
| \|Old Data 8 | \| Ex2B0080 |
| \|Old Data 9 | \| Ex2B0090 |
| \|Old Data 10 | \| EX2B0100 |
| \|Old Data 11 | \| EX2B0110 |

Note the significance of the serial numbers on the second REWND instruction and the ENDUP instruction. These serial numbers are required to inhibit the effects of an immediate action that would result in the loss of Old Data 7 and old Data 11 in the new data set. This distinction is illustrated by use of the following edit data set, in which the above-mentioned serialization does not appear:
$\begin{cases}\text { Columns 1-72 } & \text { Columns 73-80 } \\ \hdashline \text { NUM EX2B0010 } & \text { (Blank) } \\ \text { SKPTO EX2A0080 } & \text { EX2A0030 } \\ \text { REWND } & \text { EX2A0100 } \\ \text { SKPTO EX2A0070 } & \text { (Blank) } \\ \text { REWND } & \text { (Blank) } \\ \text { SKPTO EX2A0040 } & \text { (Blank) } \\ \text { SKPTO EX2A0110 } & \text { EX2A0060 } \\ \text { ENDUP } & \text { (Blank) } \\ & \end{cases}$

The following new data set is produced:


Example 3: The user wishes to update and assemble one segment of an old data set, and to create a new data set which contains the updated version of this segment. The user wishes to assemble another segment of the old data set without modification. The new data set must include this segment, as well as those segments of the old data set which the user does not wish to assemble, with the exception of one such segment, which he wishes to delete. The old data set is as follows:



A new data set is produced as shown below; records marked with an asterisk have been assembled.

Note that Example 3 will result in the production of a warning message. The operand of the last instruction in the edit data set does not correspond to a serial number in the old data set. It is speci-

| \| Columns 1-72 | \|Columns 73-80| |
| :---: | :---: |
| \|Old Data 1 | \|CPY10010 |
| \|old Data 2 | \| CPY10020 |
| \|old Data 3 | \| CPY10030 |
| \|old Data 7* | \| ASM10010 |
| IOld Data 8* | \| ASmi10020 |
| \|Old Data 9* | \| ASM10030 |
| 101d Data 10 | \|CPY20010 |
| lold Data 11 | \|CPY20020 |
| lold Data 12 | \| CPY 20030 |
| \|Old Data 13* | \|ASM20010 |
| \|O1d Data 14* | \| ASM20020 |
| \|Edit Data 1* | \| ASM20030 |
| \|Edit Data 2* | \| ASM20040 |
| \|old Data 15* | \| ASM20050 |
| \|Old Data 16* | \| ASM20060 |
| \|Edit Data 3* | \| ASM20070 |
| \|Old Data 18* | \| ASM20080 |
| \|old Data 19* | \| ASM20090 |
| \|Old Data 20 | \| CPY30010 |
| \|old Data 21 | \| CPY30020 |
| \|Old Data 22 | \|CPY30030 |

fied, however, so that the СРYTO operation will act through the last record in the old data set. A warning message will appear in the update listing.

| $\begin{aligned} & \text { 8-Bit } \\ & \text { BCD } \\ & \text { code } \end{aligned}$ | Character set Punch Combination | Decimal | Hexadecimal | Printer Graphics |
| :---: | :---: | :---: | :---: | :---: |
| 00000000 | 12,0,9,8,1 | 0 | 00 |  |
| 00000001 | 12,9,1 | 1 | 01 |  |
| 00000010 | 12,9,2 | 2 | 02 |  |
| 00000011 | 12,9,3 | 3 | 03 |  |
| 00000100 | 12,9,4 | 4 | 04 |  |
| 00000101 | 12.9 .5 | 5 | 05 |  |
| 00000110 | 12,9,6 | 6 | 06 |  |
| 00000111 | 12.9.7 | 7 | 07 |  |
| 00001000 | 12.9.8 | 8 | 08 |  |
| 00001001 | 12,9,8,1 | 9 | 09 |  |
| 00001010 | 12,9,8,2 | 10 | 0A |  |
| 00001011 | 12,9,8,3 | 11 | 0B |  |
| 00001100 | 12,9,8,4 | 12 | OC |  |
| 00001101 | 12,9,8,5 | 13 | 0D |  |
| 00001110 | 12,9,8,6 | 14 | OE |  |
| 00001111 | 12.9,8.7 | 15 | 0 F |  |
| 00010000 | 12,11,9,8,1 | 16 | 10 |  |
| 00010001 | 11,9,1 | 17 | 11 |  |
| 00010010 | 11,9,2 | 18 | 12 |  |
| 00010011 | 11,9.3 | 19 | 13 |  |
| 00010100 | 11,9.4 | 20 | 14 |  |
| 00010101 | 11,9,5 | 21 | 15 |  |
| 00010110 | 11,9,6 | 22 | 16 |  |
| 00010111 | 11,9,7 | 23 | 17 |  |
| 00011000 | 11.9.8 | 24 | 18 |  |
| 00011001 | 11,9,8,1 | 25 | 19 |  |
| 00011010 | 11,9,8,2 | 26 | 1A |  |
| 00011011 | 11,9,8,3 | 27 | 1B |  |
| 00011100 | 11,9,8,4 | 28 | 1 C |  |
| 00011101 | 11,9,8,5 | 29 | 1D |  |
| 00011110 | 11,9,8,6 | 30 | 1 E |  |
| 00011111 | 11,9,8,7 | 31 | 1 F |  |
| 00100000 | 11,0,9,8,1 | 32 | 20 |  |
| 00100001 | 0,9.1 | 33 | 21 |  |
| 00100010 | 0,9,2 | 34 | 22 |  |
| 00100011 | 0.9.3 | 35 | 23 |  |
| 00100100 | 0,9,4 | 36 | 24 |  |
| 00100101 | 0.9.5 | 37 | 25 |  |
| 00100110 | 0,9.6 | 38 | 26 |  |
| 00100111 | 0,9,7 | 39 | 27 |  |
| 00101000 | 0,9.8 | 40 | 28 |  |
| 00101001 | 0,9.8.1 | 41 | 29 |  |
| 00101010 | 0,9,8,2 | 42 | 2A |  |
| 00101011 | 0,9,8,3 | 43 | 2B |  |
| 00101100 | 0,9,8,4 | 44 | 2C |  |
| 00101101 | 0,9,8.5 | 45 | 2D |  |
| 00101110 | 0,9,8,6 | 46 | 2E |  |
| 00101111 | 0,9,8,7 | 47 | 2 F |  |
| 00110000 | 12,11,0,9,8,1 | 48 | 30 |  |
| 00110001 | 9,1 | 49 | 31 |  |
| 00110010 | 9,2 | 50 | 32 |  |


| $\begin{aligned} & \text { 8-Bit } \\ & \text { BCD } \\ & \text { code } \end{aligned}$ | Character Set Punch combination | Decimal | Hexadecimal | Printer Graphics |
| :---: | :---: | :---: | :---: | :---: |
| 00110011 | 9.3 | 51 | 33 |  |
| 00110100 | 9.4 | 52 | 34 |  |
| 00110101 | 9.5 | 53 | 35 |  |
| 00110110 | 9.6 | 54 | 36 |  |
| 00110111 | 9.7 | 55 | 37 |  |
| 00111000 | 9,8 | 56 | 38 |  |
| 00111001 | 9,8,1 | 57 | 39 |  |
| 00111010 | 9.8.2 | 58 | 3A |  |
| 00111011 | 9.8.3 | 59 | 3B |  |
| 00111100 | 9,8,4 | 60 | 3 C |  |
| 00111101 | 9.8.5 | 61 | 3D |  |
| 00111110 | 9,8,6 | 62 | 3E |  |
| 00111111 | 9,8,7 | 63 | 3 F |  |
| 01000000 |  | 64 | 40 | blank |
| 01000001 | 12.0.9.1 | 65 | 41 |  |
| 01000010 | 12,0,9,2 | 66 | 42 |  |
| 01000011 | 12,0,9.3 | 67 | 43 |  |
| 01000100 | 12.0.9.4 | 68 | 44 |  |
| 01000101 | 12,0,9,5 | 69 | 45 |  |
| 01000110 | 12,0,9,6 | 70 | 46 |  |
| 01000111 | 12,0,9,7 | 71 | 47 |  |
| 01001000 | 12,0,9,8 | 72 | 48 |  |
| 01001001 | 12,8,1 | 73 | 49 |  |
| 01001010 | 12,8,2 | 74 | 4A |  |
| 01001011 | 12,8.3 | 75 | 4 B | - (period) |
| 01001100 | 12.8.4 | 76 | 4 C |  |
| 01001101 | 12,8,5 | 77 | 4D | ( |
| 01001110 | 12,8,6 | 78 | 4E | + |
| 01001111 | 12,8,7 | 79 | 4 F |  |
| 01010000 | 12 | 80 | 50 | $\varepsilon$ |
| 01010001 | 12.11.9.1 | 81 | 51 |  |
| 01010010 | 12,11,9,2 | 82 | 52 |  |
| 01010011 | 12,11,9,3 | 83 | 53 |  |
| 01010100 | 12.11.9.4 | 84 | 54 |  |
| 01010101 | 12,11,9,5 | 85 | 55 |  |
| 01010110 | 12.11.9.6 | 86 | 56 |  |
| 01010111 | 12,11,9,7 | 87 | 57 |  |
| 01011000 | 12,11,9,8 | 88 | 58 |  |
| 01011001 | 11,8,1 | 89 | 59 |  |
| 01011010 | 11.8.2 | 90 | 5A |  |
| 01011011 | 11,8,3 | 91 | 5B | \$ |
| 01011100 | 11,8,4 | 92 | 5c | * |
| 01011101 | 11,8,5 | 93 | 5D | ) |
| 01011110 | 11,8,6 | 94 | 5 E |  |
| 01011111 | 11,8,7 | 95 | 5 F |  |
| 01100000 | 11 | 96 | 60 | - |
| 01100001 | 0.1 | 97 | 61 | / |
| 01100010 | 11,0,9,2 | 98 | 62 |  |
| 01100011 | 11,0,9.3 | 99 | 63 |  |
| 01100100 | 11,0,9,4 | 100 | 64 |  |
| 01100101 | 11,0,9,5 | 101 | 65 |  |
| 01100110 | 11,0,9,6 | 102 | 66 |  |
| 01100111 | 11,0,9,7 | 103 | 67 |  |
| 01101000 | 11.0.9.8 | 104 | 68 |  |
| 01101001 | 0,8,1 | 105 | 69 |  |
| 01101010 | 12.11 | 106 | 6A |  |
| 01101011 | 0,8,3 | 107 | 6B | , (comma) |


| $\begin{aligned} & 8-\text { Bit } \\ & \text { BCD } \\ & \text { code } \end{aligned}$ | Character Set Punch Combination | Decimal | Hexadecimal | Printer Graphics |
| :---: | :---: | :---: | :---: | :---: |
| 01101100 | 0,8,4 | 108 | 6C | \% |
| 01101101 | 0,8,5 | 109 | 6D |  |
| 01101110 | 0,8,6 | 110 | 6 E |  |
| 01101111 | 0,8,7 | 111 | 6 F |  |
| 01110000 | 12,11,0 | 112 | 70 |  |
| 01110001 | 12,11,0,9.1 | 113 | 71 |  |
| 01110010 | 12,11,0,9,2 | 114 | 72 |  |
| 01110011 | 12,11,0,9,3 | 115 | 73 |  |
| 01110100 | 12,11,0,9.4 | 116 | 74 |  |
| 01110101 | 12.11,0,9.5 | 117 | 75 |  |
| 01110110 | 12,11,0,9,6 | 116 | 76 |  |
| 01110111 | 12,11,0,9,7 | 119 | 77 |  |
| 01111000 | 12,11,0,9,8 | 120 | 78 |  |
| 01111001 | 8.1 | 121 | 79 |  |
| 01111010 | 8,2 | 122 | 7A |  |
| 01111011 | 8.3 | 123 | 7 B | \# |
| 01111100 | 8,4 | 124 | 7 C | จ |
| 01111101 | 8,5 | 125 | 7D | - (single quote) |
| 01111110 | 8,6 | 126 | 7 E | $=$ |
| 01111111 | 8,7 | 127 | 7 F |  |
| 10000000 | 12,0,8,1 | 128 | 80 |  |
| 10000001 | 12,0,1 | 129 | 81 |  |
| 10000010 | 12,0,2 | 130 | 82 |  |
| 10000011 | 12,0,3 | 131 | 83 |  |
| 10000100 | 12.0.4 | 132 | 84 |  |
| 10000101 | 12.0 .5 | 133 | 85 |  |
| 10000110 | 12,0,6 | 134 | 86 |  |
| 10000111 | 12,0,7 | 135 | 87 |  |
| 10001000 | 12,0,8 | 136 | 88 |  |
| 10001001 | 12,0,9 | 137 | 89 |  |
| 10001010 | 12,0,8,2 | 138 | 8A |  |
| 10001011 | 12,0,8,3 | 139 | 8B |  |
| 10001100 | 12,0,8,4 | 140 | 8 C |  |
| 10001101 | 12,0,8,5 | 141 | 8D |  |
| 10001110 | 12,0,8,6 | 142 | 8 E |  |
| 10001111 | 12,0,8.7 | 143 | 8 F |  |
| 10010000 | 12,11,8,1 | 144 | 90 |  |
| 10010001 | 12,11.1 | 145 | 91 |  |
| 10010010 | 12,11,2 | 146 | 92 |  |
| 10010011 | 12,11.3 | 147 | 93 |  |
| 10010100 | 12,11,4 | 148 | 94 |  |
| 10010101 | 12,11,5 | 149 | 95 |  |
| 10010110 | 12,11,6 | 150 | 96 |  |
| 10010111 | 12.11.7 | 151 | 97 |  |
| 10011000 | 12,11,8 | 152 | 98 |  |
| 10011001 | 12,11,9 | 153 | 99 |  |
| 10011010 | 12,11,8,2 | 154 | 9A |  |
| 10011011 | 12,11.8.3 | 155 | 9 B |  |
| 10011100 | 12,11,8,4 | 156 | 9 C |  |
| 10011101 | 12,11,8,5 | 157 | 9 D |  |
| 10011110 | 12,11,8,6 | 158 | 9 E |  |
| 10011111 | 12,11,8,7 | 159 | 9 F |  |
| 10100000 | 11,0,8,1 | 160 | A0 |  |
| 10100001 | 11,0,1 | 161 | A1 |  |
| 10100010 | 11,0,2 | 162 | A2 |  |
| 10100011 | 11,0,3 | 163 | A3 |  |
| 10100100 | 11.0,4 | 164 | A 4 |  |



| 8-Bit | Character set |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| $B C D$ | Punch |  | Hexa- | Printer |
| code | Combination | Decimal | decimal | Graphics |
| 11011110 | 12,11,9,8,6 | 222 | DE |  |
| 11011111 | $12,11,9,8,7$ | 223 | DF |  |
| 11100000 | 0,8,2 | 224 | E0 |  |
| 11100001 | 11,0,9.1 | 225 | E1 |  |
| 11100010 | 0,2 | 226 | E2 | S |
| 11100011 | 0,3 | 227 | E3 | T |
| 11100100 | 0.4 | 228 | E4 | U |
| 11100101 | 0,5 | 229 | E5 | V |
| 11100110 | 0.6 | 230 | E6 | w |
| 11100111 | 0.7 | 231 | E7 | X |
| 11101000 | 0,8 | 232 | E8 | Y |
| 11101001 | 0,9 | 233 | E9 | Z |
| 11101010 | $11,0,9,8,2$ | 234 | EA |  |
| 11101011 | 11, 0,9,8,3 | 235 | EB |  |
| 11101100 | 11,0,9,8,4 | 236 | EC |  |
| 11101101 | 11,0,9,8,5 | 237 | ED |  |
| 11101110 | 11,0,9,8,6 | 238 | EE |  |
| 11101111 | 11,0,9,8,7 | 239 | EF |  |
| 11110000 | 0 | 240 | F0 | 0 |
| 11110001 | 1 | 241 | F1 | 1 |
| 11110010 | 2 | 242 | F2 | 2 |
| 11110011 | 3 | 243 | F3 | 3 |
| 11110100 | 4 | 244 | F4 | 4 |
| 11110101 | 5 | 245 | F5 | 5 |
| 11110110 | 6 | 246 | F6 | 6 |
| 11110111 | 7 | 247 | F7 | 7 |
| 11111000 | 8 | 248 | F8 | 8 |
| 11111001 | 9 | 249 | F9 | 9 |
| 11111010 | 12,11,0,9,8,2 | 250 | FA |  |
| 11111011 | 12,11,0,9,8,3 | 251 | FB |  |
| 11111100 | $12,11,0,9,8,4$ | 252 | FC |  |
| 11111101 | $12,11,0,9,8,5$ | 253 | FD |  |
| 11111110 | $12,11,0,9,8,6$ | 254 | FE |  |
| 11111111 | $12,11,0,9,8,7$ | 255 | FF |  |

The table provides direct conversion of decimal and hexadecimal numbers in the following ranges:


Decimal numbers (0000-4095) are given within the table. The first two characters (high-order) of hexadecimal numbers (000-FFF) are given in the left-hand column of the table; the third character (x) is arranged across the top of each part of the table.

To find the decimal equivalent of the hexadecimal number 0C9, look for 0C in the left-hand column, and across that row under the column for $\mathrm{x}=9$. The decimal number is 0201.

To convert from decimal to hexadecimal, look up the decimal number within the table and read the hexadecimal number by a combination of the hexadecimal characters in the left-hand column, and the value for $x$ at
the top of the column containing the decimal number. For example, the decimal number 123 has the hexadecimal equivalent of 07B; the decimal number 1478 has the hexadecimal equivalent of 5 C 6 .

For numbers outside the range of the table, add the following values to the table.

| Hexadecimal | Decimal |
| :---: | :---: |
| 1000 | 4096 |
| 2000 | 8192 |
| 3000 | 12288 |
| 5000 | 16384 |
| 6000 | 20480 |
| 7000 | 24576 |
| 8000 | 28672 |
| 9000 | 36868 |
| A000 | 40960 |
| B000 | 45056 |
| C000 | 49152 |
| D000 | 53248 |
| F000 | 57344 |
|  | 61440 |


|  | $\mathrm{x}=0$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | $F$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00x | 0000 | 0007 | 0002 | 0003 | 0004 | 0005 | 0006 | 0007 | 00008 | 00009 | 0010 | 0011 | 0012 | 0013 | ÔO 14 | 0015 |
| 01x | 0016 | 0017 | 0018 | 0019 | 0020 | 0021 | 0022 | 0023 | 0024 | 0025 | 0026 | 0027 | 0028 | 0029 | 0030 | 0031 |
| 02x | 0032 | 0033 | 0034 | 0035 | 0036 | 0037 | 0038 | 0039 | 0040 | 0041 | 0042 | 0043 | 0044 | 0045 | 0046 | 0047 |
| 03x | 0048 | 0049 | 0050 | 0051 | 0052 | 0053 | 0054 | 0055 | 0056 | 0057 | 0058 | 0059 | 0060 | 0061 | 0062 | 000. |
| 04x | 0064 | 0065 | 0066 | 0067 | 0068 | 0069 | 0070 | 0071 | 0072 | 0073 | 0074 | 0075 | 0076 | 0077 | 0078 | 007 |
| 05x | 0080 | 0081 | 0082 | 0083 | 0084 | 0085 | 0086 | 0087 | 0088 | 0089 | 0090 | 0091 | 0092 | 0093 | 0094 | 0095 |
| 06x | 0096 | 0097 | 0098 | 0099 | 0100 | 0101 | 0102 | 0103 | 0104 | 0105 | 0106 | 0107 | 0108 | 0109 | 0110 | 0111 |
| 07x | 0112 | 0113 | 0114 | 0115 | 0116 | 0117 | 0118 | 0119 | 0120 | 0121 | 0122 | 0123 | 0124 | 0125 | 0126 | 012 ? |
| 08x | 0128 | 0129 | 0130 | 0131 | 0132 | 0133 | 0134 | 0135 | 0136 | 0137 | 0138 | 0139 | 0140 | 0141 | 0142 | 0143 |
| 09x | 0144 | 0145 | 0146 | 0147 | 0148 | 0149 | 0150 | 0151 | 0152 | 0153 | 0154 | 0155 | 0156 | 0157 | 0158 | 0153 |
| 0Ax | 0160 | 0161 | 0162 | 0163 | 0164 | 0165 | 0166 | 0167 | 0168 | 0169 | 0170 | 0171 | 0172 | 0173 | 0174 | 0175 |
| 08x | 0176 | 0177 | 0178 | 0179 | 0180 | 0181 | 0182 | 0183 | 0184 | 0185 | 0186 | 0187 | 0188 | 0189 | 0190 | 0191 |
| 0 Cx | 0192 | 0193 | 0194 | 0195 | 0196 | 0197 | 0198 | 0199 | 0200 | 0201 | 0202 | 0203 | 0204 | 0205 | 0206 | 0207 |
| 0Dx | 0208 | 0209 | 0210 | 0211 | 0212 | 0213 | 0214 | 0215 | 0216 | 0217 | 0218 | 0219 | 0220 | 0221 | 0222 | 0223 |
| UEX | 0224 | 0225 | 0226 | 0227 | 0228 | 0229 | 0230 | 0231 | 0232 | 0233 | 0234 | 0235 | 0236 | 0237 | 0238 | 0239 |
| 0Fx | 0240 | 0241 | 0242 | 0243 | 0244 | 0245 | 0246 | 0247 | 0248 | 0249 | 0250 | 0251 | 0252 | 0253 | 0254 | 0255 |
| 10x | 0256 | 0257 | 0258 | 0259 | 0260 | 0261 | 0262 | 0263 | 0264 | 0265 | 0266 | 0267 | 0268 | 0269 | 0270 | 0271 |
| 11x | 0272 | 0273 | 0274 | 0275 | 0276 | 0277 | 0278 | 0279 | 0280 | 0281 | 0282 | 0283 | 0284 | 0285 | 0286 | 0287 |
| 12x | 0288 | 0289 | 0290 | 0291 | 0292 | 0293 | 0294 | 0295 | 0296 | 0297 | 0298 | 0299 | 0300 | 0301 | 0302 | 0303 |
| 13x | 0304 | 0305 | 0306 | 0307 | 0308 | 0309 | 0310 | 0311 | 0312 | 0313 | 0314 | 0315 | 0316 | 0317 | 0318 | 0319 |
| 14 x | 0320 | 0321 | 0322 | 0323 | 0324 | 0325 | 0326 | 0327 | 0328 | 0329 | 0330 | 0331 | 0332 | 0333 | 0334 | 0335 |
| 15x | 0336 | 0337 | 0338 | 0339 | 0340 | 0341 | 0342 | 0343 | 0344 | 0345 | 0346 | 0347 | 0348 | 0349 | 0350 | 0351 |
| 16x | 0352 | 0353 | 0354 | 0355 | 0356 | 0357 | 0358 | 0359 | 0360 | 0361 | 0362 | 0363 | 0364 | 0365 | 0366 | 0367 |
| 17x | 0368 | 0369 | 0370 | 0371 | 0372 | 0373 | 0374 | 0375 | 0376 | 0377 | 0378 | 0379 | 0380 | 0381 | 0382 | 0383 |
| 18x | 0384 | 0385 | 0386 | 0387 | 0388 | 0389 | 0390 | 0391 | 0392 | 0393 | 0394 | 0395 | 0396 | 0397 | 0398 | 0399 |
| 19x | 0400 | 0401 | 0402 | 0403 | 0404 | 0405 | 0406 | 0407 | 0408 | 0409 | 0410 | 0411 | 0412 | 0413 | 0414 | 0415 |
| 1 Ax | 0416 | 0417 | 0418 | 0419 | 0420 | 0421 | 0422 | 0423 | 0424 | 0425 | 0426 | 0427 | 0428 | 0429 | 0430 | 0431 |
| 1 Bx | 0432 | 0433 | 0434 | 0435 | 0436 | 0437 | 0438 | 0439 | 0440 | 0441 | 0442 | 0443 | 0444 | 0445 | 0446 | 0447 |
| 1Cx | 0448 | 0449 | 0450 | 0451 | 0452 | 0453 | 0454 | 0455 | 0456 | 0457 | 0458 | 0459 | 0460 | 0461 | 0462 | 006? |
| 1Dx | 0464 | 0465 | 0466 | 0467 | 0468 | 0469 | 0470 | 0471 | 3472 | 0473 | 0474 | 0475 | 0476 | 0477 | 0478 | 0479 |
| 1 Ex | 0480 | 0481 | 0482 | 0483 | 0484 | 0485 | 0486 | 0487 | 0488 | 0489 | 0490 | 0491 | 0492 | 0493 | 0494 | 045 |
| 1Fx | 0496 | 0497 | 0498 | 0499 | 0500 | 0501 | 0502 | 0503 | 0504 | 0505 | 0506 | 0507 | 0508 | 0509 | 0510 | 0511 |
| 20x | 0512 | 0513 | 0514 | 0515 | 0516 | 0517 | 0518 | 0519 | 0520 | 0521 | 0522 | 0523 | 0524 | 0525 | 0526 | 0527 |
| 21x | 0528 | 0529 | 0530 | 0531 | 0532 | 0533 | 0534 | 0535 | 0536 | 0537 | 0538 | 0539 | 0540 | 0541 | 0542 | 0543 |
| 22x | 0544 | 0545 | 0546 | 0547 | 0548 | 0549 | 0550 | 0551 | 0552 | 0553 | 0554 | 0555 | 0556 | 0557 | 0558 | 0559 |
| 23x | 0560 | 0561 | 0562 | 0563 | 0564 | 0565 | 0566 | 0567 | 0568 | 0569 | 0570 | 0571 | 0572 | 0573 | 0574 | 0575 |
| 24x | 0576 | 0577 | 0578 | 0579 | 0580 | 0581 | 0582 | 0583 | 0584 | 0585 | 0586 | 0587 | 0588 | 0589 | 0590 | 0591 |
| 25x | 0592 | 0593 | 0594 | 0595 | 0596 | 0597 | 0598 | 0599 | 0600 | 0601 | 0602 | 0603 | 0604 | 0605 | 0606 | 0607 |
| 26x | 0608 | 0609 | 0610 | 0611 | 061 ? | 0613 | 0614 | 0615 | 0616 | 0617 | 0618 | 0619 | 0620 | 0621 | 0622 | 0623 |
| 27x | 0624 | 0625 | 0626 | 0627 | 0628 | 0629 | 0630 | 0631 | 0632 | 0633 | 0634 | 0635 | 0636 | 0637 | 0638 | 0639 |
| 28x | 0640 | 0641 | 0642 | 0643 | 0644 | 0645 | 0646 | 0647 | 0648 | 0649 | 0650 | 0651 | 0652 | 0653 | 0654 | 0655 |
| 29x | 0656 | 0657 | 0658 | 0659 | 0660 | 0661 | 0662 | 0663 | 0664 | 0665 | 0666 | 0667 | 0668 | 0669 | 0670 | 0671 |
| 2Ax | 0672 | 0673 | 0674 | 0675 | 0676 | 0677 | 0678 | 0679 | 0680 | 0681 | 0682 | 0683 | 0684 | 0685 | 0686 | 0687 |
| 2Bx | 0688 | 0689 | 0690 | 0691 | 0692 | 0693 | 0694 | 0695 | 0696 | 0697 | 0698 | 0699 | 0700 | 0701 | 0702 | 0703 |
| 2 Cx | 0704 | 0705 | 0706 | 0707 | 0708 | 0709 | 0710 | 0711 | 0712 | 0713 | 0714 | 0715 | 0716 | 0717 | 0718 | 0719 |
| 2Dx | 0720 | 0721 | 0722 | 0723 | 0724 | 0725 | 0726 | 0727 | 0728 | 0729 | 0730 | 0731 | 0732 | 0733 | 0734 | 0735 |
| 2Ex | 0736 | 0737 | 0738 | 0739 | 0740 | 0741 | 0742 | 0743 | 0744 | 0745 | 0746 | 0747 | 0748 | 0749 | 0750 | 0751 |
| 2 Fx | 0752 | 0753 | 0754 | 0755 | 0756 | 0757 | 0758 | 0759 | 0760 | 0761 | 0762 | 0763 | 0764 | 0765 | 0766 | 0767 |
| 30x | 0768 | 0769 | 0770 | 0771 | 0772 | 0773 | 0774 | 0775 | 0776 | 0777 | 0778 | 0779 | 0780 | 0781 | 0782 | 0783 |
| $31 \times$ | 0784 | 0785 | 0786 | 0787 | 0788 | 0789 | 0790 | 0791 | 0792 | 0793 | 0794 | 0795 | 0796 | 0797 | 0798 | 0799 |
| 32x | 0800 | 0801 | 0802 | 0803 | 0804 | 0805 | 0806 | 0807 | 0808 | 0809 | 0810 | 0811 | 0812 | 0813 | 0814 | 0815 |
| 33x | 0816 | 0817 | 0818 | 0819 | 0820 | 0821 | 0822 | 0823 | 0824 | 0825 | 0826 | 0827 | 0828 | 0829 | 0830 | 0831 |
| 34x | 0832 | 0833 | 0834 | 0835 | 0836 | 0837 | 0838 | 0839 | 0840 | 0841 | 0842 | 0843 | 0844 | 0845 | 0846 | 0847 |
| 35x | 0848 | 0849 | 0850 | 0851 | 0852 | 0853 | 0854 | 0855 | 0856 | 0857 | 0858 | 0859 | 0860 | 0861 | 0862 | 0863 |
| 36x | 0864 | 0865 | 0866 | 0867 | 0868 | 0869 | 0870 | 0871 | 0872 | 0873 | 0874 | 0875 | 0876 | 0877 | 0878 | 0879 |
| $37 x$ | 0880 | 0881 | 0882 | 0883 | 0884 | 0885 | 0886 | 0887 | 0888 | 0889 | 0890 | 0891 | 0892 | 0893 | 0894 | 0895 |
| 38x | 0896 | 0897 | 0898 | 0899 | 0900 | 0901 | 0902 | 0903 | 0904 | 0905 | 0906 | 0907 | 0908 | 0909 | 0910 | 0911 |
| 39x | 0912 | 0913 | 0914 | 0915 | 0916 | 0917 | 0918 | 0919 | 0920 | 0921 | 0922 | 0923 | 0924 | 0925 | 0926 | 0927 |
| 3Ax | 0928 | 0929 | 0930 | 0931 | 0932 | 0933 | 0934 | 0935 | 0936 | 0937 | 0938 | 0939 | 0940 | 0941 | 0942 | 0943 |
| 3Bx | 0944 | 0945 | 0946 | 0947 | 0948 | 0949 | 0950 | 0951 | 0952 | 0953 | 0954 | 0955 | 0956 | 0957 | 0958 | 0959 |
| 3Cx | 0960 | 0961 | 0962 | 0963 | 0964 | 0965 | 0966 | 0967 | 0968 | 0969 | 0970 | 0971 | 0972 | 0973 | 0974 | 0975 |
| 3Dx | 0976 | 0977 | 0978 | 0979 | 0980 | 0981 | 0982 | 0983 | 0984 | 0985 | 0986 | 0987 | 0988 | 0989 | 0990 | 0991 |
| 3Ex | 0992 | 0993 | 0994 | 0995 | 0996 | 0997 | 0998 | 0999 | 1000 | 1001 | 1002 | 1003 | 1004 | 1005 | 1006 | 1007 |
| 3 Fx | 1008 | 1009 | 1010 | 1011 | 1012 | 1013 | 1014 | 1015 | 1016 | 1017 | 1018 | 1019 | 1020 | 1021 | 1022 | 1023 |


|  | $\mathrm{x}=0$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | $F$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 40x | 1024 | 1025 | 1026 | 1027 | 1028 | 1029 | 1030 | 1031 | 1032 | 1033 | 1034 | 1035 | 1036 | 1037 | 1038 | 1039 |
| 41 x | 1040 | 1041 | 1042 | 1043 | 1044 | 1045 | 1046 | 1047 | 1048 | 1049 | 1050 | 1051 | 1052 | 1053 | 1054 | 1055 |
| 42 x | 1056 | 1057 | 1058 | 1059 | 1060 | 1061 | 1062 | 1063 | 1064 | 1065 | 1066 | 1067 | 1068 | 1069 | 1070 | 1071 |
| 43x | 1072 | 1073 | 1074 | 1075 | 1076 | 1077 | 1078 | 1079 | 1080 | 1081 | 1082 | 1083 | 1084 | 1085 | 1086 | 1087 |
| 44 x | 1088 | 1089 | 1090 | 1091 | 1092 | 1093 | 1094 | 1095 | 1096 | 1097 | 1098 | 1099 | 1100 | 1101 | 1102 | 1103 |
| 45x | 1104 | 1105 | 1106 | 1107 | 1108 | 1109 | 1110 | 1111 | 1112 | 1113 | 1114 | 1115 | 1116 | 1117 | 1118 | 1119 |
| 46x | 1120 | 1121 | 1122 | 1123 | 1124 | 1125 | 1126 | 1127 | 1128 | 1129 | 1130 | 1131 | 1132 | 1133 | 1134 | 1135 |
| 47x | 1136 | 1137 | 1138 | 1139 | 1140 | 1141 | 1142 | 1143 | 1144 | 1145 | 1146 | 1147 | 1148 | 1149 | 1150 | 1151 |
| 48x | 1152 | 1153 | 1154 | 1155 | 1156 | 1157 | 1158 | 1159 | 1160 | 1161 | 1162 | 1163 | 1164 | 1165 | 1166 | 1167 |
| 49x | 1168 | 1169 | 1170 | 1171 | 1172 | 1173 | 1174 | 1175 | 1176 | 1177 | 1178 | 1179 | 1180 | 1181 | 1182 | 1183 |
| 4Ax | 1184 | 1185 | 1186 | 1187 | 1188 | 1189 | 1190 | 1191 | 1192 | 1193 | 1194 | 1195 | 1196 | 1197 | 1198 | 1199 |
| 4 Bx | 1200 | 1201 | 1202 | 1203 | 1204 | 1205 | 1206 | 1207 | 1208 | 1209 | 1210 | 1211 | 1212 | 1213 | 1214 | 1215 |
| 4 Cx | 1216 | 1217 | 1218 | 1219 | 1220 | 1221 | 1222 | 1223 | 1224 | 1225 | 1226 | 1227 | 1228 | 1229 | 1230 | 1231 |
| 4Dx | 1232 | 1233 | 1234 | 1235 | 1236 | 1237 | 1238 | 1239 | 1240 | 1241 | 1242 | 1243 | 1244 | 1245 | 1246 | 1247 |
| 4Ex | 1248 | 1249 | 1250 | 1251 | 1252 | 1253 | 1254 | 1255 | 1256 | 1257 | 1258 | 1259 | 1260 | 1261 | 1262 | 1263 |
| 4Fx | 1264 | 1265 | 1266 | 1267 | 1268 | 1269 | 1270 | 1271 | 1272 | 1273 | 1274 | 1275 | 1276 | 1277 | 1278 | 1279 |
| 50x | 1280 | 1281 | 1282 | 1283 | 1284 | 1285 | 1286 | 1287 | 1288 | 1289 | 1290 | 1291 | 1292 | 1293 | 1294 | 1295 |
| 51x | 1296 | 1297 | 1298 | 1299 | 1300 | 1301 | 1302 | 1303 | 1304 | 1305 | 1306 | 1307 | 1308 | 1309 | 1310 | 1311 |
| 52x | 1312 | 1313 | 1314 | 1315 | 1316 | 1317 | 1318 | 1319 | 1320 | 1321 | 1322 | 1323 | 1324 | 1325 | 1326 | 1327 |
| 53x | 1328 | 1329 | 1330 | 1331 | 1332 | 1333 | 1334 | 1335 | 1336 | 1337 | 1338 | 1339 | 1340 | 1341 | 1342 | 1343 |
| 54x | 1344 | 1345 | 1346 | 1347 | 1348 | 1349 | 1350 | 1351 | 1352 | 1353 | 1354 | 1355 | 1356 | 1357 | 1358 | 1359 |
| 55x | 1360 | 1361 | 1362 | 1363 | 1364 | 1365 | 1366 | 1367 | 1368 | 1369 | 1370 | 1371 | 1372 | 1373 | 1374 | 1375 |
| 56x | 1376 | 1377 | 1378 | 1379 | 1380 | 1381 | 1382 | 1383 | 1384 | 1385 | 1386 | 1387 | 1388 | 1389 | 1390 | 1391 |
| 57x | 1392 | 1393 | 1394 | 1395 | 1396 | 1397 | 1398 | 1399 | 1400 | 1401 | 1402 | 1403 | 1404 | 1405 | 1406 | 1407 |
| 58x | 1408 | 1409 | 1410 | 1411 | 1412 | 1413 | 1414 | 1415 | 1416 | 1417 | 1418 | 1419 | 1420 | 1421 | 1422 | 1423 |
| 59x | 1424 | 1425 | 1426 | 1427 | 1428 | 1429 | 1430 | 1431 | 1432 | 1433 | 1434 | 1435 | 1436 | 1437 | 1438 | 1439 |
| 5Ax | 1440 | 1441 | 1442 | 1443 | 1444 | 1445 | 1446 | 1447 | 1448 | 1449 | 1450 | 1451 | 1452 | 1453 | 1454 | 1455 |
| 5Bx | 1456 | 1457 | 1458 | 1459 | 1460 | 1461 | 1462 | 1463 | 1464 | 1465 | 1466 | 1467 | 1468 | 1469 | 1470 | 1471 |
| 5Cx | 1472 | 1473 | 1474 | 1475 | 1476 | 1477 | 1478 | 1479 | 1480 | 1481 | 1482 | 1483 | 1484 | 1485 | 1486 | 1487 |
| 5Dx | 1488 | 1489 | 1490 | 1491 | 1492 | 1493 | 1494 | 1495 | 1496 | 1497 | 1498 | 1499 | 1500 | 1501 | 1502 | 1503 |
| 5Ex | 1504 | 1505 | 1506 | 1507 | 1508 | 1509 | 1510 | 1511 | 1512 | 1513 | 1514 | 1515 | 1516 | 1517 | 1518 | 1519 |
| 5Fx | 1520 | 1521 | 1522 | 1523 | 1524 | 1525 | 1526 | 1527 | 1528 | 1529 | 1530 | 1531 | 1532 | 1533 | 1534 | 1535 |
| 60 x | 1536 | 1537 | 1538 | 1539 | 1540 | 1541 | 1542 | 1543 | 1544 | 1545 | 1546 | 1547 | 1548 | 1549 | 1550 | iS51 |
| $61 \times$ | 1552 | 1553 | 1554 | 1555 | 1556 | 1557 | 1558 | 1559 | 1560 | 1561 | 1562 | 1563 | 1564 | 1565 | 1566 | 1567 |
| 62x | 1568 | 1569 | 1570 | 1571 | 1572 | 1573 | 1574 | 1575 | 1576 | 1577 | 1578 | 1579 | 1580 | 1581 | 1582 | 1583 |
| 63x | 1584 | 1585 | 1586 | 1587 | 1588 | 1589 | 1590 | 1591 | 1592 | 1593 | 1594 | 1595 | 1596 | 1597 | 1598 | 1599 |
| 64 x | 1600 | 1601 | 1602 | 1603 | 1604 | 1605 | 1606 | 1607 | 1608 | 1609 | 1610 | 1611 | 1612 | 1613 | 1614 | 1615 |
| $65 x$ | 1616 | 1617 | 1618 | 1619 | 1620 | 1621 | 1622 | 1623 | 1624 | 1625 | 1626 | 1627 | 1628 | 1629 | 1630 | 1631 |
| 66x | 1632 | 1633 | 1634 | 1635 | 1636 | 1637 | 1638 | 1639 | 1640 | 1641 | 1642 | 1643 | 1644 | 1645 | 1646 | 1647 |
| 67x | 1648 | 1649 | 1650 | 1651 | 1652 | 1653 | 1654 | 1655 | 1656 | 1657 | 1658 | 1659 | 1660 | 1661 | 1662 | 1663 |
| 68 x | 1664 | 1665 | 1666 | 1667 | 1668 | 1669 | 1670 | 1671 | 1672 | 1673 | 1674 | 1675 | 1676 | 1677 | 1678 | 1679 |
| 69x | 1680 | 1681 | 1682 | 1683 | 1684 | 1685 | 1686 | 1687 | 1688 | 1689 | 1690 | 1691 | 1692 | 1693 | 1694 | 1695 |
| 6Ax | 1696 | 1697 | 1698 | 1699 | 1700 | 1701 | 1702 | 1703 | 1704 | 1705 | 1706 | 1707 | 1708 | 1709 | 1710 | 1711 |
| 6 Bx | 1712 | 1713 | 1714 | 1715 | 1716 | 1717 | 1718 | 1719 | 1720 | 1721 | 1722 | 1723 | 1724 | 1725 | 1726 | 1727 |
| 6 Cx | 1728 | 1729 | 1730 | 1731 | 1732 | 1733 | 1734 | 1735 | 1736 | 1737 | 1738 | 1739 | 1740 | 1741 | 1742 | 1743 |
| 6Dx | 1744 | 1745 | 1746 | 1747 | 1748 | 1749 | 1750 | 1751 | 1752 | 1753 | 1754 | 1755 | 1756 | 1757 | 1758 | 1759 |
| 6 Ex | 1760 | 1761 | 1762 | 1763 | 1764 | 1765 | 1766 | 1767 | 1768 | 1769 | 1770 | 1771 | 1772 | 1773 | 1774 | 1775 |
| 6 Fx | 1776 | 1777 | 1778 | 1779 | 1780 | 1781 | 1782 | 1783 | 1784 | 1785 | 1786 | 1787 | 1788 | 1789 | 1790 | 1791 |
| 70x | 1792 | 1793 | 1794 | 1795 | 1796 | 1797 | 1798 | 1799 | 1800 | 1801 | 1802 | 1803 | 1804 | 1805 | 1806 | 1807 |
| 71x | 1808 | 1809 | 1810 | 1811 | 1812 | 1813 | 1814 | 1815 | 1816 | 1817 | 1818 | 1819 | 1820 | 1821 | 1822 | 1823 |
| 72x | 1824 | 1825 | 1826 | 1827 | 1828 | 1829 | 1830 | 1831 | 1832 | 1833 | 1834 | 1835 | 1836 | 1837 | 1838 | 1839 |
| 73x | 1840 | 1841 | 1842 | 1843 | 1844 | 1845 | 1846 | 1847 | 1848 | 1849 | 1850 | 1851 | 1852 | 1853 | 1854 | 1855 |
| 74 x | 1856 | 1857 | 1858 | 1859 | 1860 | 1861 | 1862 | 1863 | 1864 | 1865 | 1866 | 1867 | 1868 | 1869 | 1870 | 1871 |
| 75x | 1872 | 1873 | 1874 | 1875 | 1876 | 1877 | 1878 | 1879 | 1880 | 1881 | 1882 | 1883 | 1884 | 1885 | 1886 | 1887 |
| 76x | 1888 | 1889 | 1890 | 1891 | 1892 | 1893 | 1894 | 1895 | 1896 | 1897 | 1898 | 1899 | 1900 | 1901 | 1902 | 1903 |
| 77x | 1904 | 1905 | 1906 | 1907 | 1908 | 1909 | 1910 | 1911 | 1912 | 1913 | 1914 | 1915 | 1916 | 1917 | 1918 | 1919 |
| 78x | 1920 | 1921 | 1922 | 1923 | 1924 | 1925 | 1926 | 1927 | 1928 | 1929 | 1930 | 1931 | 1932 | 1933 | 1934 | 1935 |
| 79x | 1936 | 1937 | 1938 | 1939 | 1940 | 1941 | 1942 | 1943 | 1944 | 1945 | 1946 | 1947 | 1948 | 1949 | 1950 | 1951 |
| 7Ax | 1952 | 1953 | 1954 | 1955 | 1956 | 1957 | 1958 | 1959 | 1960 | 1961 | 1962 | 1963 | 1964 | 1965 | 1966 | 1967 |
| 7 Bx | 1968 | 1969 | 1970 | 1971 | 1972 | 1973 | 1974 | 1975 | 1976 | 1977 | 1978 | 1979 | 1980 | 198: | 1982 | 1983 |
| 7 Cx | 1984 | 1985 | 1986 | 1987 | 1988 | 1989 | 1990 | 1991 | 1992 | 1993 | 1994 | 1995 | 1996 | 1997 | 1998 | 1999 |
| 7Dx | 2000 | 2001 | 2002 | 2003 | 2004 | 2005 | 2006 | 2007 | 2008 | 2009 | 2010 | 2011 | 2012 | 2013 | 2014 | 2015 |
| 7Ex | 2016 | 2017 | 2018 | 2019 | 2020 | 2021 | 2022 | 2023 | 2024 | 2025 | 2026 | 2027 | 2028 | 2029 | 2030 | 2031 |
| 7Fx | 2032 | 2033 | 2034 | 2035 | 2036 | 2037 | 2038 | 2039 | 2040 | 2041 | 2042 | 2043 | 2044 | 2045 | 2046 | 2047 |


|  | $\mathbf{x}=0$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | c | D | E | F |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 80x | 2048 | 2049 | 2050 | 2051 | 2052 | 2053 | 2054 | 2055 | 2056 | 2057 | 2058 | 2059 | 2060 | 2061 | 2062 | 2063 |
| 81x | 2064 | 2065 | 2066 | 2067 | 2068 | 2069 | 2070 | 2071 | 2072 | 2073 | 2074 | 2075 | 2076 | 2077 | 2078 | 2079 |
| 82x | 2080 | 2081 | 2082 | 2083 | 2084 | 2085 | 2086 | 2087 | 2088 | 2089 | 2090 | 2091 | 2092 | 2093 | 2094 | 2095 |
| 83x | 2096 | 2097 | 2098 | 2099 | 2100 | 2101 | 2102 | 2103 | 2104 | 2105 | 2106 | 2107 | 2108 | 2109 | 2110 | 2111 |
| 84 x | 2112 | 2113 | 2114 | 2115 | 2116 | 2117 | 2118 | 2119 | 2120 | 2121 | 2122 | 2123 | 2124 | 2125 | 2126 | 2127 |
| 85 x | 2128 | 2129 | 2130 | 2131 | 2132 | 2133 | 2134 | 2135 | 2136 | 2137 | 2138 | 2139 | 2140 | 2141 | 2142 | 2143 |
| 86x | 2144 | 2145 | 2146 | 2147 | 2148 | 2149 | 2150 | 2151 | 2152 | 2153 | 2154 | 2155 | 2156 | 2157 | 2158 | 2159 |
| 87 x | 2160 | 2161 | 2162 | 2163 | 2164 | 2165 | 2166 | 2167 | 2168 | 2169 | 2170 | 2171 | 2172 | 2173 | 2174 | 2175 |
| 88x | 2176 | 2177 | 2178 | 2179 | 2180 | 2181 | 2182 | 2183 | 2184 | 2185 | 2186 | 2187 | 2188 | 2189 | 2190 | 2191 |
| 89x | 2192 | 2193 | 2194 | 2195 | 2196 | 2197 | 2198 | 2199 | 2200 | 2201 | 2202 | 2203 | 2204 | 2205 | 2206 | 2207 |
| 8Ax | 2208 | 2209 | 2210 | 2211 | 2212 | 2213 | 2214 | 2215 | 2216 | 2217 | 2218 | 2219 | 2220 | 2221 | 2222 | 2223 |
| 8 Bx | 2224 | 2225 | 2226 | 2227 | 2228 | 2229 | 2230 | 2231 | 2232 | 2233 | 2234 | 2235 | 2236 | 2237 | 2238 | 2239 |
| 8 Cx | 2240 | 2241 | 2242 | 2243 | 2244 | 2245 | 2246 | 2247 | 2248 | 2249 | 2250 | 2251 | 2252 | 2253 | 2254 | 2255 |
| 8Dx | 2256 | 2257 | 2258 | 2259 | 2260 | 2261 | 2262 | 2263 | 2264 | 2265 | 2266 | 2267 | 2268 | 2269 | 2270 | 2271 |
| 8Ex | 2272 | 2273 | 2274 | 2275 | 2276 | 2277 | 2278 | 2279 | 2280 | 2281 | 2282 | 2283 | 2284 | 2285 | 2286 | 2287 |
| 8Fx | 2288 | 2289 | 2290 | 2291 | 2292 | 2293 | 2294 | 2295 | 2296 | 2297 | 2298 | 2299 | 2300 | 2301 | 2302 | 2303 |
| 90x | 2304 | 2305 | 2306 | 2307 | 2308 | 2309 | 2310 | 23:1 | 2312 | 2313 | 2314 | 2315 | 2316 | 2317 | 2318 | 2319 |
| 91x | 2320 | 2321 | 2322 | 2323 | 2324 | 2325 | 2326 | 2327 | 2328 | 2329 | 2330 | 2331 | 2332 | 2333 | 2334 | 2335 |
| 92x | 2336 | 2337 | 2338 | 2339 | 2340 | 2341 | 2342 | 2343 | 2344 | 2345 | 2346 | 2347 | 2348 | 2349 | 2350 | 2351 |
| 93x | 2352 | 2353 | 2354 | 2355 | 2356 | 2357 | 2358 | 2359 | 2360 | 2361 | 2362 | 2363 | 2364 | 2365 | 2366 | 2367 |
| 94x | 2368 | 2369 | 2370 | 2371 | 2372 | 2373 | 2374 | 2375 | 2376 | 2377 | 2378 | 2379 | 2380 | 2381 | 2382 | 2383 |
| 95x | 2384 | 2385 | 2386 | 2387 | 2388 | 2389 | 2390 | 2391 | 2392 | 2393 | 2394 | 2395 | 2396 | 2397 | 2398 | 2399 |
| 96x | 2400 | 2401 | 2402 | 2403 | 2404 | 2405 | 2406 | 2407 | 2408 | 2409 | 2410 | 2411 | 2412 | 2413 | 2414 | 2415 |
| 97x | 2416 | 2417 | 2418 | 2419 | 2420 | 2421 | 2422 | 2423 | 2424 | 2425 | 2426 | 2427 | 2428 | 2429 | 2430 | 2431 |
| 98x | 2432 | 2433 | 2434 | 2435 | 2436 | 2437 | 2438 | 2439 | 2440 | 2441 | 2442 | 2443 | 2444 | 2445 | 2446 | 2447 |
| 99x | 2448 | 2449 | 2450 | 2451 | 2452 | 2453 | 2454 | 2455 | 2456 | 2457 | 2458 | 2459 | 2460 | 2461 | 2462 | 2463 |
| 9Ax | 2464 | 2465 | 2466 | 2467 | 2468 | 2469 | 2470 | 2471 | 2472 | 2473 | 2474 | 2475 | 2476 | 2477 | 2478 | 2479 |
| 9Bx | 2480 | 2481 | 2482 | 2483 | 2484 | 2485 | 2486 | 2487 | 2488 | 2489 | 2490 | 2491 | 2492 | 2493 | 2494 | 2495 |
| 9 Cx | 2496 | 2497 | 2498 | 2499 | 2500 | 2501 | 2502 | 2503 | 2504 | 2505 | 2506 | 2507 | 2508 | 2509 | 2510 | 2511 |
| 9Dx | 2512 | 2513 | 2514 | 2515 | 2516 | 2517 | 2518 | 2519 | 2520 | 2521 | 2522 | 2523 | 2524 | 2525 | 2526 | 2527 |
| 9Ex | 2528 | 2529 | 2530 | 2531 | 2532 | 2533 | 2534 | 2535 | 2536 | 2537 | 2538 | 2539 | 2540 | 2541 | 2542 | 2543 |
| 9Fx | 2544 | 2545 | 2546 | 2547 | 2548 | 2549 | 2550 | 2551 | 2552 | 2553 | 2554 | 2555 | 2556 | 2557 | 2558 | 2559 |
| A0x | 2560 | 2561 | 2562 | 2563 | 2564 | 2565 | 2566 | 2567 | 2568 | 2569 | 2570 | 2571 | 2572 | 2573 | 2574 | 2575 |
| A1x | 2576 | 2577 | 2578 | 2579 | 2580 | 2581 | 2582 | 2583 | 2584 | 2585 | 2586 | 2587 | 2588 | 2589 | 2590 | 2591 |
| A2x | 2592 | 2593 | 2594 | 2595 | 2596 | 2597 | 2598 | 2599 | 2600 | 2601 | 2602 | 2603 | 2604 | 2605 | 2606 | 2607 |
| A3x | 2608 | 2609 | 2610 | 2611 | 2612 | 2613 | 2614 | 2615 | 2616 | 2617 | 2618 | 2619 | 2620 | 2621 | 2622 | 2623 |
| ${ }^{\text {A } 4 x}$ | 2624 | 2625 | 2626 | 2627 | 2628 | 2629 | 2630 | 2631 | 2632 | 2633 | 2634 | 2635 | 2636 | 2637 | 2638 | 2639 |
| A5 $x$ | 2640 | 2641 | 2642 | 2643 | 2644 | 2645 | 2646 | 2647 | 2648 | 2649 | 2650 | 2651 | 2652 | 2653 | 2654 | 2655 |
| A6x | 2656 | 2657 | 2658 | 2659 | 2660 | 2661 | 2662 | 2663 | 2664 | 2665 | 2666 | 2667 | 2668 | 2669 | 2670 | 2671 |
| A7x | 2672 | 2673 | 2674 | 2675 | 2676 | 2677 | 2678 | 2679 | 2680 | 2681 | 2682 | 2683 | 2684 | 2685 | 2686 | 2687 |
| A8x | 2688 | 2689 | 2630 | 2691 | 2692 | 2693 | 2694 | 2695 | 2696 | 2697 | 2698 | 2699 | 2700 | 2701 | 2702 | 2703 |
| A9x | 2704 | 2705 | 2706 | 2707 | 2708 | 2709 | 2710 | 2711 | 2712 | 2713 | 2714 | 2715 | 2716 | 2717 | 2718 | 2719 |
| AAx | 2720 | 2721 | 2722 | 2723 | 2724 | 2725 | 2726 | 2727 | 2728 | 2729 | 2730 | 2731 | 2732 | 2733 | 2734 | 2735 |
| ABx | 2736 | 2737 | 2738 | 2739 | 2740 | 2741 | 2742 | 2743 | 2744 | 2745 | 2746 | 2747 | 2748 | 2749 | 2750 | 2751 |
| ACx | 2752 | 2753 | 2754 | 2755 | 2756 | 2757 | 2758 | 2759 | 2760 | 2761 | 2762 | 2763 | 276" | ${ }^{7} 765$ | 2766 | 2767 |
| ADx | 2768 | 2769 | 2770 | 2771 | 2772 | 2773 | 2774 | 2775 | 2776 | 2777 | 2778 | 2779 | 2780 | 2781 | 2782 | 2783 |
| AEx | 2784 | 2785 | 2786 | 2787 | 2788 | 2789 | 2790 | 2791 | 2792 | 2793 | 2794 | 2795 | 2796 | 2797 | 2798 | 2799 |
| AFx | 2800 | 2801 | 2802 | 2803 | 2804 | 2805 | 2806 | 2807 | 2808 | 2809 | 2810 | 2811 | 2812 | 2813 | 2814 | 2815 |
| B0x | 2816 | 2817 | 2818 | 2819 | 2820 | 2821 | 2822 | 2823 | 2824 | 2825 | 2826 | 2827 | 2828 | 2829 | 2830 | 2831 |
| B1x | 2832 | 2833 | 2834 | 2835 | 2836 | 2837 | 2838 | 2839 | 2840 | 2841 | 2842 | 2843 | 2844 | 2845 | 2846 | 2847 |
| B2x | 2848 | 2849 | 2850 | 2851 | 2852 | 2853 | 2854 | 2855 | 2856 | 2857 | 2858 | 2859 | 2860 | 2861 | 2862 | 2863 |
| B3x | 2864 | 2865 | 2866 | 2867 | 2868 | 2869 | 2870 | 2871 | 2872 | 2873 | 2874 | 2875 | 2876 | 2877 | 2878 | 2879 |
| B4x | 2880 | 2881 | 2882 | 2883 | 2884 | 2885 | 2886 | 2887 | 2888 | 2889 | 2890 | 2891 | 2892 | 2893 | 2894 | 2895 |
| B5x | 2896 | 2897 | 2898 | 2899 | 2900 | 2901 | 2902 | 2903 | 2904 | 2905 | 2906 | 2907 | 2908 | 2909 | 2910 | 2911 |
| B6x | 2912 | 2913 | 2914 | 2915 | 2916 | 2917 | 2918 | 2919 | 2920 | 2921 | 2922 | 2923 | 2924 | 2925 | 2926 | 2927 |
| B7x | 2928 | 2929 | 2930 | 2931 | 2932 | 2933 | 2934 | 2935 | 2936 | 2937 | 2938 | 2939 | 2940 | 2941 | 2942 | 2943 |
| B8x | 2944 | 2945 | 2946 | 2947 | 2948 | 2949 | 2950 | 2951 | 2952 | 2953 | 2954 | 2955 | 2956 | 2957 | 2958 | 2959 |
| B9x | 2960 | 2961 | 2962 | 2963 | 2964 | 2965 | 2966 | 2967 | 2968 | 2969 | 2970 | 2971 | 2972 | 2973 | 2974 | 2975 |
| BAx | 2976 | 2977 | 2978 | 2979 | 2980 | 2981 | 2982 | 2983 | 2984 | 2985 | 2986 | 2987 | 2988 | 2989 | 2990 | 2991 |
| BBx | 2992 | 2993 | 2994 | 2995 | 2996 | 2997 | 2998 | 2999 | 3000 | 3001 | 3002 | 3003 | 3004 | 3005 | 3006 | 3007 |
| BCx | 3008 | 3009 | 3010 | 3011 | 3012 | 3013 | 3014 | 3015 | 3016 | 3017 | 3018 | 3019 | 3020 | 3021 | 302. | 3023 |
| BDx | 3024 | 3025 | 3026 | 3027 | 3028 | 3029 | 3030 | 3031 | 3032 | 3033 | 3034 | 3035 | 3036 | 3037 | 3038 | 3039 |
| BEx | 3040 | 3041 | 3042 | 3043 | 3044 | 3045 | 3046 | 3047 | 3048 | 3049 | 3050 | 3051 | 3052 | 3053 | 3054 | 3055 |
| BFx | 3056 | 3057 | 3058 | 3059 | 3060 | 3061 | 3062 | 3063 | 3064 | 3065 | 3066 | 3067 | 3068 | 3069 | 3070 | 3071 |


№tes for Appendix $C$ :

1. R1, R2, and R3 are absolute terms that specify qeneral or floating-point reqisters. The qeneral register numbers are 0 throuqh 15; floating-point reqister numbers are $0,2,4$, and 6. In $B C$ and $B C R$ machine instructions, $R, ~ s p e c i f i e s ~ a ~ 4-b i t ~ m a s k . ~$
2. D1 and D2 are absolute expressions that specify displacements. A value of 0 throuqh 4095 may be specified.
3. B1 and B2 are absolute terms that specify base reqisters. Register numbers are 0 throuqh 15.
4. $x 2$ is an absolute term that specifies an index register. Reqister numbers are 0 throuqh 15.
5. I and I2 are absolute expressions that provide immediate data. The value of the expression may be 0 through 255.
6. S1 and 52 are absolute or relocatable expressions that specify an address.
7. RR, RS, and SI instruction fields that are crossed out in the machine formats are not examined during instruction execution. The fields are not written in the symbolic operand, but are assembled as binary zeros.

This appendix contains an alphabetical listing of the mnemonic operation codes of all the machine instructions that can be represented in assembler lanquaqe. The column headinas in the list and the information each column provides are as follows:

Mnemonic Code: This column gives the mnemonic operation code for the machine instruction.

Instruction: This column contains the name of the instruction associated with the mnemonic.

Machine Code: This column contains the hexadecimal equivalent of the actual machine operation code.

Basic Machine Format: This column qives the basic machine format of the instruction: RR, RX, RS, or SI.

Operand Field Format: This column shows the symbolic format of the operand field for the particular memonic.



Form C28-6.811-1, paqe revised 5/20/68, by TNL N33-8543


## APPENDIX E: ASSEMBLER INSTRUCTIONS

| Operation | Name Field | Operand Field |
| :---: | :---: | :---: |
| - AGO | Sequence symbol optional | A sequence symbol |
| \| AIF | Sequence symbol optional | A logical expression immediately followed by a sequence symbol |
| \| ANOP | Sequence symbol required | Must be blank |
| CCW | Symbol optional | Four operands |
| - CNOP | Must be blank | Two absolute expressions, separated by a comma |
| \| COM | Symbol optional | Must be blank |
| CSECT | Symbol optional ${ }^{1}$ | Must be blank |
| DC | Symbol optional | One operand |
| \| DROP | Must be blank | One operand |
| - DS | Symbol optional | One operand |
| 1 DSECT | Symbol required | Must be blank |
| \| EJECT | Must be blank | Must be blank |
| \| END | Must be blank | A relocatable expression or blank |
| ENTRY | Must be blank | One operand |
| I EQU | Symbol required | An absolute or relocatable expression |
| EXTRN | Must be blank | One operand |
| ICTI. | Must be blank | The decimal value 1 or 25 |
| LTORG | Symbol optional | Must be blank |
| ORG | Must be blank | A relocatable expression or blank |
| 1 PRINT | Must be blank | One or two operands |
| - REPro | Must be blank | Must be blank |
| SETA | Variable symbol required | An arithmetic expression |
| SPACE | Must be blank | A decimal self-defining term or blank |
| START | Symbol optional | A self-defining term or blank |
| TITLE ${ }^{2}$ | A special symbol (1 to 4 characters) or not present | One to 62 characters, enclosed in single quotes |
| USING | Must be blank | An absolute or relocatable expression followed by an absolute term |
| ${ }^{1}$ A symbol during th ${ }^{2}$ See Sect | required if an unnamed START or assembly. <br> 5 for the description of the nam | ECT instruction has already been used entry. |




#### Abstract

With certain exceptions, the IBM System/ 360 Model 44 Proqramming System Assembler Lanquage is a selected subset of the languages available in the IBM System/360 proqramming support systems designed for the Models $30,40,50,65$, and 75 -specifically, System/360 Operating System (0S/360), System/360 Disk Operating System (DOS/360), and System/360 Tape Operating System (TOS/360). This appendix uses three lists to describe this subset. The first of these lists comprises those features of the Model 44 Programming System Assembler Lanquage that are not within the subset (i.e., those features peculiar to the Model 44 Assembler). The second and third lists comprise, respectively, those features of the other system/360 assembler languaqes that are not included in the Model 44 Programming System Assembler Lanquage, and those that are included subject to the limitations noted.


1. Features of the IBM System/360 Model 44 Programminq System Assembler Langaaqe that are not supported by the other System/360 assembler languages are, as follows:
a. Update instructions
b. Named common control sections
c. Implicit definition of SETA symbols
d. Mnemonic operation codes for the following machine instructions:
(1) Chanqe Priority Mask (CHPM)
(2) Load PSW Special (LPSX)
(3) Read Direct Word (RDDW)
(4) Write Direct Word (WRDW)
2. Features of the IBM System/360 Operatinq System Assembler Lanquaqe (and, in some cases, of some or all of the other System/360 programminq support system assembler lanquages) that are not supported by the Model 44 Proqramming System Assembler Lanquaqe are, as follows:
a. Biñày self-afefiñing teiulu
b. Continuation cards
c. COPY instruction
d. ISEQ instruction
e. Macro instructions
f. Mnemonic operation codes for machine instructions in the storage-to-storage (SS) format, or for any of the following machine instructions:
(1) Convert to Binary (CVB)
(2) Convert to Decimal (CVD)
(3) Read Direct (RDD)
(4) Write Direct (WRD)
q. PUNCH instruction
3. Features of the IBM System/360 Operating System Assembler Lanquaqe (and, in some cases, of some or all of.the other System/360 proqramming support system assembler lanquaqes) only selected subsets of which are supported by the Model 44 Programming System Assembler Lanquage are, as follows:


| 1 FEATURE | 1 LIMITATIONS |
| :---: | :---: |
|  | + |
| IICTL instruction | Ione operand only (thel |
|  | \|begin column specifi-| |
|  | \|cation); operand must| |
|  | \|assume a value of 1 or |
|  | \|25. | |
|  |  |
| \| Literal pool | IThe positioning of thel |
|  | \|literal pool must bel |
|  | \|assiqned by the pro-1 |
|  | \|qrammer, and a literal| |
|  | \|pool must be assiqned| |
|  | \|following the last| |
|  | \|occurrence of a literal| |
|  | \|in any qiven control| |
|  | \|section. | |
|  |  |
| ```\|Machine instruction operands``` | IOperands R1, R2, B1.1 |
|  | \|B2, and X2 (seel |
|  | \|Appendix C) must be| |
|  | \|written as absolute| |
|  | Iterms; multiterm abso-1 |
|  | Ilute expressions arel |
|  | \|not permitted for thesel |
|  | loperands. |
|  |  |
| \|PRINT instruction | I Two operands only (thel |
|  | \|GEN/NOGEN option is notl |
|  | (supported). |
|  |  |
| \|REPRO instruction | IColumns 1 through 721 |
|  | Ionly are reproduced. |
|  | --m--- |
| ITITLE instruction | \|Maximum of 62 charac-1 |
|  | \|ters in operand field| |
|  | \| (exclusive of delimit-1 |
|  | \|ing sinqle quotes). | |
|  | +--1 |
| \|USING instruction| | loperand $\underline{r}$ must be an \| |
| 1 | \|absolute term; only onel |
| 1 | \|operand r (i.e.f only| |
| 1 | \|one qeneral reqister) | |
| 1 | Imay be specified in al |
| 1 | \|single USING| |
| 1 | Iinstruction. |
|  |  |

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Substitute the following pages for the pages currently in your copy of the publication:

> 35 and 36
> 67 and 68
> 69 and 70
> 73 and 74

All changes and additions are indicated by a vertical line to the left of the affected portion of the text. Figures that were changed have a bullet (•) to the left of the figure caption.

File this cover letter at the back of the publication. It will then serve as a record of the changes received and incorporated.

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[^0]:    ๑ 1966 by International Business Machines Corporation

[^1]:    * Discussed in Section 3.

[^2]:    * Not permitted for character and hexadecimal constants.

